

Beyond real: Investigating the role of complex numbers in self-testing

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Abstract

We investigate complex self-testing, a generalization of standard self-testing that accounts for quantum strategies whose statistics is indistinguishable from their complex conjugate's. We show that many structural results from standard self-testing extend to the complex setting, including lifting of common assumptions. Our main result is an operator-algebraic characterization: complex self-testing is equivalent to uniqueness of the real parts of higher moments, leading to a basis-independent formulation in terms of real C* algebras. This leads to a classification of non-local strategies, and a tight boundary where standard self-testing does not apply and complex self-testing is necessary. We further construct a strategy involving quaternions, establishing the first standard self-test for genuinely complex strategy. Our work clarifies the structure of complex self-testing and highlights the subtle role of complex numbers in bipartite Bell non-locality.

Keywords: Bell Scenario, entanglement, self-testing, device-independence, C* algebra, complex strategy, real strategy.

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Contents

1	Introduction	3
2	Preliminaries and notion	4
2.1	Bell scenarios	4
2.2	Complex dilation and complex self-testing	5
3	Basic properties, and lifting assumptions in complex self-testing	8
4	An operator-algebraic characterization	12
4.1	Complex self-testing implies a unique real part of moments	12
4.2	Unique real part of moments implies complex self-testing	13
4.3	The operator-algebraic formulation	15
5	Realness of quantum strategies	17
6	A quaternion middle ground possibility	23
6.1	A self-test involving quaternions	23
6.2	Generating quaternions by few projections	27

1 Introduction

Self-testing is a powerful concept in quantum information theory that enables the certification of quantum states and measurements solely from observed correlations. This idea originates from Bell non-locality [Bel64, CHSH69]. In the 1980s, studies identified the maximal quantum violation of the CHSH inequality, showing that it is uniquely attained by a specific entangled state and measurement setup [Tsi87, PR92]. This uniqueness serves as the foundation for the notion of self-testing, formalized by Mayers and Yao [MY04]. Since its inception, self-testing has evolved into an active and expanding field of research. The utility of self-testing extends broadly and it lies behind applications across quantum information science [CGJV19, BCM⁺18, MY04, VV14, JNV⁺20]. A comprehensive overview can be found in [ŠB20].

It is a folklore fact that complex strategies cannot be self-tested, at least not in the standard sense [ŠB20, Section 3.7.1]. To see this, consider a strategy $S = (|\psi\rangle, \{E_{xa}\}, \{F_{yb}\})$ for a (bipartite) Bell scenario; that is, Alice performs $\{E_{xa}\}$, and Bob performs $\{F_{yb}\}$, on their respective parts of the shared state $|\psi\rangle$. The resulting correlation satisfies

$$p(a, b|x, y) = \langle \psi | E_{xa} \otimes F_{yb} | \psi \rangle = \overline{\langle \bar{\psi} | E_{xa} \otimes F_{yb} | \bar{\psi} \rangle}.$$

That is, S and $\bar{S} := (|\bar{\psi}\rangle, \{\overline{E_{xa}}\}, \{\overline{F_{yb}}\})$ always give rise to a same correlation p . Since there is generally no local unitary mapping between S and \bar{S} ,¹ the standard self-testing framework, which certifies uniqueness up to local unitaries, is an ill fit for such cases. This is especially relevant in multipartite Bell scenarios where quantum states cannot always be represented by real numbers [ŠBR⁺23, BJCA⁺24].

To address this, the notion of *complex self-testing* [MM11] was introduced, allowing any combination of a strategy and its complex conjugate. This idea has been explored in several works, first in the bipartite scenario [MM11, APVW16, BŠCA18, JMS20], and also in multipartite scenarios [ŠBR⁺23, BJCA⁺24]. It is worth noting that, certifying complex measurements is tied to deep foundational problems in quantum information theory [RTW⁺21]. However, complex self-testing remains relatively underdeveloped in comparison to the rich theory of (standard) self-testing. Many questions regarding its structural properties and limits are yet to be answered.

In this work, we undertake a systematic study of complex self-testing. Our main contributions are as follows:

- We show that many fundamental properties of standard self-testing carry over to the complex setting. This generalizes the results of removing common assumptions like projectivity and full-rankness from [BCK⁺25].
- We provide an operator-algebraic characterization of complex self-testing. Specifically, we show that complex self-testing is equivalent to uniqueness of the real part of higher moments across all strategies reproducing the same correlation. This leads to a natural basis-independent definition of complex self-testing, in terms of unique real states on real C* algebras—contrasting with the (complex) C* algebra framework used for standard self-testing [PSZZ24].

¹A simple example is a strategy contains all three Pauli measurements: there is no unitary map that takes $(\sigma_X, \sigma_Y, \sigma_Z)$ to $(\sigma_X, \bar{\sigma}_Y = -\sigma_Y, \sigma_Z)$

- As a consequence, we prove that complex self-testing reduces to standard self-testing whenever the canonical strategy has only real-valued moments. This observation prompts a deeper investigation into the “realness” of quantum strategies and leads us to identify a fundamental boundary: strategies with nonreal higher moments cannot be certified via standard self-testing.
- Finally, we explore an intermediate regime between real-representable strategies and those with real higher moments. We construct a new self-test that relies on quaternions, demonstrating that nontrivial (standard) self-tests exist in this middle ground. A technical by-product of our construction is a tight lower bound on the number of projections required to generate the quaternion matrix algebra. This gives the minimal scenario for this type of self-tested strategies to exist, and might be of independent interest from an operator-algebraic point of view.

Through this work, we aim to provide a coherent and comprehensive foundation for complex self-testing, clarify its relationship to standard self-testing, and open new directions for understanding the algebraic structure underlying non-locality.

2 Preliminaries and notion

Throughout this paper all Hilbert spaces (denoted by \mathcal{H} , with subscripts indicating the party they belong to) are assumed to be over the complex field and finite-dimensional, unless specified otherwise. The set of linear operators on Hilbert space is denoted by $L(\mathcal{H})$. The identity operator of a d -dimensional Hilbert space is denoted by Id , with subscripts indicating the party. For a matrix a , let a^* , \bar{a} , a^\top denote its adjoint, complex conjugate, and transpose, respectively. A (pure) state (namely, a unit vector) in a Hilbert space $\mathcal{H} = \mathbb{C}^n$ is denoted by $|\psi\rangle$ and we use $|\bar{\psi}\rangle$ to denote its complex conjugate. Real, complex, and quaternion numbers are denoted by \mathbb{R} , \mathbb{C} , and \mathbb{H} , respectively.

2.1 Bell scenarios

In a standard bipartite Bell scenario [Bel64, BCP⁺14], non-communicating players Alice and Bob are modelled as performing local measurements on their respective part of their shared state. Their behaviour then is described as *quantum strategies*:

$$S = \left(|\psi\rangle_{AB} \in \mathcal{H}_A \otimes \mathcal{H}_B, \{E_{xa} : x \in \mathcal{I}_A, a \in \mathcal{O}_A\} \subset L(\mathcal{H}_A), \{F_{yb} : y \in \mathcal{I}_B, b \in \mathcal{O}_B\} \subset L(\mathcal{H}_B) \right).$$

In each round of the interaction, a verifier samples inputs $x \in \mathcal{I}_A$ and $y \in \mathcal{I}_B$ from finite question sets $\mathcal{I}_{A,B}$ and sends them to Alice and Bob respectively. In response, players produce outputs $a \in \mathcal{O}_A$ and $b \in \mathcal{O}_B$ from finite answer sets $\mathcal{O}_{A,B}$, based on local quantum measurements $\{E_{xa}\}, \{F_{yb}\}$ applied to their respective subsystems. Therefore, the statistics, also called the *correlation* of S , is given by the conditional probability distribution $p(a, b|x, y) = \langle \psi | E_{xa} \otimes F_{yb} | \psi \rangle$, which can be estimated through repeated executions of the protocol. The set of all possible correlations generated by finite dimensional tensor product quantum strategies is denoted by C_q .

In [BCK⁺25], basic properties about non-local strategies were introduced.

Definition 2.1. A strategy $S = (|\psi\rangle \in \mathcal{H}_A \otimes \mathcal{H}_B, \{E_{xa}\}, \{F_{yb}\})$ is

- support-preserving if

$$[\Pi_A, E_{xa}] = [\Pi_B, F_{yb}] = 0,$$

hold for all x, y, a, b , where Π_A (resp. Π_B) is the projection onto the support of $\text{Tr}_B |\psi\rangle\langle\psi|$ (resp. $\text{Tr}_A |\psi\rangle\langle\psi|$);

- full-rank if $|\psi\rangle$ has full Schmidt rank;
- 0-projective (or ‘projective on the state’) if

$$\langle\psi|(\text{Id}_A - E_{xa})E_{xa} \otimes \text{Id}_B|\psi\rangle = \langle\psi|\text{Id}_A \otimes (\text{Id}_B - F_{yb})F_{yb}|\psi\rangle = 0,$$

hold for all x, y, a, b ;

- projective if $E_{xa}^2 = E_{xa}, F_{yb}^2 = F_{yb}$ holds for all a, b, x, y .

Clearly, full-rank implies support-preserving, and projective implies 0-projective.

The following property of support-preservingness will be useful in our proofs.

Lemma 2.2 (Lemmas 4.3 & 4.4 of [PSZZ24]; see also Lemma 3.3 of [BCK⁺25] for the approximate version). A strategy $(|\psi\rangle, \{E_{xa}\}, \{F_{yb}\})$ is support-preserving if and only if there exist operators $\hat{E}_{xa}, \hat{F}_{yb}$ such that $E_{xa} \otimes \text{Id} |\psi\rangle = \text{Id} \otimes \hat{E}_{xa} |\psi\rangle$ and $\text{Id} \otimes F_{yb} |\psi\rangle = \hat{F}_{yb} \otimes \text{Id} |\psi\rangle$ for all x, y, a, b .

A strategy S as above is *irreducible* if neither $\{E_{xa}\}$ nor $\{F_{yb}\}$ have a non-trivial proper closed invariant subspace. If S is finite-dimensional (in the sense that $\dim \mathcal{H}_A, \dim \mathcal{H}_B < \infty$), this is equivalent to $\{E_{xa}\}$ and $\{F_{yb}\}$ generating $L(\mathcal{H}_A)$ and $L(\mathcal{H}_B)$ as complex algebras.

For the sake of simplicity, we denote words (products of operators) of length k by $E_{\vec{x}\vec{a}} := E_{x_k a_k} E_{x_{k-1} a_{k-1}} \cdots E_{x_1 a_1}$, $\tilde{E}_{\vec{x}\vec{a}} := \tilde{E}_{x_k a_k} \tilde{E}_{x_{k-1} a_{k-1}} \cdots \tilde{E}_{x_1 a_1}$, where $\vec{x} := (x_k, \dots, x_1), \vec{a} := (a_k, \dots, a_1)$. Similarly for Bob’s operators, $F_{\vec{y}\vec{b}} := F_{y_\ell b_\ell} F_{y_{\ell-1} b_{\ell-1}} \cdots F_{y_1 b_1}$, $\tilde{F}_{\vec{y}\vec{b}} := \tilde{F}_{y_\ell b_\ell} \tilde{F}_{y_{\ell-1} b_{\ell-1}} \cdots \tilde{F}_{y_1 b_1}$ denote words of length ℓ . One may view the correlation $\langle\psi|E_{xa} \otimes F_{yb}|\psi\rangle$ of S as the first moments of the joint distribution of Alice and Bob. In line with this view, we call $\langle\psi|E_{\vec{x}\vec{a}} \otimes F_{\vec{y}\vec{b}}|\psi\rangle$ the *higher moments* of the strategy S .

2.2 Complex dilation and complex self-testing

Self-testing aims to establish a correspondence between the *canonical* strategy and the *physical* strategy. The canonical strategy $\tilde{S} = (|\tilde{\psi}\rangle_{AB}, \{\tilde{E}_{xa}\}, \{\tilde{F}_{yb}\})$ is the specification or the blueprint to be compared with, while the physical strategy $S = (|\psi\rangle_{AB}, \{E_{xa}\}, \{F_{yb}\})$ is performed by the players. The notion of local dilation, introduced in [MPS24, PSZZ24] and now standard in the literature, describes such a correspondence (more precisely, a partial order), incorporating undetectable auxiliary resource and change of the frame of reference.

Definition 2.3 (Local dilation). A strategy $\tilde{S} = (|\tilde{\psi}\rangle, \{\tilde{E}_{xa}\}, \{\tilde{F}_{yb}\})$ is a local dilation of a strategy $S = (|\psi\rangle, \{E_{xa}\}, \{F_{yb}\})$ (denoted $S \hookrightarrow \tilde{S}$) if there exist a local isometry $U = U_A \otimes U_B$ with

$$\begin{aligned} U_A &: \mathcal{H}_A \rightarrow \mathcal{H}_{\tilde{A}} \otimes \mathcal{H}_{\hat{A}}, \\ U_B &: \mathcal{H}_B \rightarrow \mathcal{H}_{\tilde{B}} \otimes \mathcal{H}_{\hat{B}} \end{aligned}$$

and an auxiliary state $|j\rangle$ such that

$$\begin{aligned} U|\psi\rangle_{AB} &= |\tilde{\psi}\rangle_{\tilde{A}\tilde{B}} |j\rangle_{\hat{A}\hat{B}} \\ U(E_{xa} \otimes \text{Id}_B)|\psi\rangle_{AB} &= (\tilde{E}_{xa} \otimes \text{Id}_{\tilde{B}} |\tilde{\psi}\rangle_{\tilde{A}\tilde{B}}) |j\rangle_{\hat{A}\hat{B}} \\ U(\text{Id}_A \otimes F_{yb})|\psi\rangle_{AB} &= (\text{Id}_{\tilde{A}} \otimes \tilde{F}_{yb} |\tilde{\psi}\rangle_{\tilde{A}\tilde{B}}) |j\rangle_{\hat{A}\hat{B}} \end{aligned}$$

hold for all a, b, x, y .

In a *complex* self-test, the physical strategy is expected to be an arbitrary combination of \tilde{S} and its complex conjugate. We formulate this with a *complex local dilation*, the ‘complex’ analogue of a local dilation. The idea is to introduce an additional Hilbert space $\mathcal{H}_{A'} \otimes \mathcal{H}_{B'}$ on which the devices perform measurements on an entangled state to concurrently employ the canonical strategy or its complex conjugate. Since the additional measurements acting on $\mathcal{H}_{A'}$ and $\mathcal{H}_{B'}$ has binary outcomes, without loss of generality we take $\mathcal{H}_{A'} \cong \mathcal{H}_{B'} \cong \mathbb{C}^2$, and the state in $\mathcal{H}(A') \otimes \mathcal{H}(B')$ can take the form $\alpha|00\rangle + \beta|11\rangle$. Also notice that the real coefficients α, β can be absorbed to the auxiliary states. We then define complex local dilation as follows.

Definition 2.4 (Complex local dilation). A strategy $\tilde{S} = (|\tilde{\psi}\rangle, \{\tilde{E}_{xa}\}, \{\tilde{F}_{yb}\})$ is a complex local dilation of a strategy $S = (|\psi\rangle, \{E_{xa}\}, \{F_{yb}\})$ (denoted $S \hookrightarrow_{\mathbb{C}} \tilde{S}$) if there exists a local isometry $U = U_A \otimes U_B$ with

$$\begin{aligned} U_A &: \mathcal{H}_A \rightarrow \mathcal{H}_{\tilde{A}} \otimes \mathcal{H}_{\hat{A}} \otimes \mathcal{H}_{A'}, \\ U_B &: \mathcal{H}_B \rightarrow \mathcal{H}_{\tilde{B}} \otimes \mathcal{H}_{\hat{B}} \otimes \mathcal{H}_{B'} \end{aligned}$$

such that

$$U|\psi\rangle_{AB} = |\tilde{\psi}\rangle_{\tilde{A}\tilde{B}} |j_0\rangle_{\hat{A}\hat{B}} |00\rangle_{A'B'} + |\overline{\tilde{\psi}}\rangle_{\tilde{A}\tilde{B}} |j_1\rangle_{\hat{A}\hat{B}} |11\rangle_{A'B'}, \quad (1)$$

$$U(E_{xa} \otimes \text{Id}_B)|\psi\rangle_{AB} = (\tilde{E}_{xa} \otimes \text{Id}_{\tilde{B}} |\tilde{\psi}\rangle_{\tilde{A}\tilde{B}}) |j_0\rangle_{\hat{A}\hat{B}} |00\rangle_{A'B'} + (\overline{\tilde{E}_{xa}} \otimes \text{Id}_{\tilde{B}} |\overline{\tilde{\psi}}\rangle_{\tilde{A}\tilde{B}}) |j_1\rangle_{\hat{A}\hat{B}} |11\rangle_{A'B'}, \quad (2)$$

$$U(\text{Id}_A \otimes F_{yb})|\psi\rangle_{AB} = (\text{Id}_{\tilde{A}} \otimes \tilde{F}_{yb} |\tilde{\psi}\rangle_{\tilde{A}\tilde{B}}) |j_0\rangle_{\hat{A}\hat{B}} |00\rangle_{A'B'} + (\text{Id}_{\tilde{A}} \otimes \overline{\tilde{F}_{yb}} |\overline{\tilde{\psi}}\rangle_{\tilde{A}\tilde{B}}) |j_1\rangle_{\hat{A}\hat{B}} |11\rangle_{A'B'} \quad (3)$$

hold for all a, b, x, y , where $|j_0\rangle$ and $|j_1\rangle$ are (not necessarily orthogonal) subnormalized states satisfying $\langle j_0|j_0\rangle + \langle j_1|j_1\rangle = 1$.

Clearly, if \tilde{S} is already represented by real matrices, $|00\rangle_{A'B'}$ and $|11\rangle_{A'B'}$ can be absorbed into the auxiliary state, in which case the complex local dilation degenerates to a (standard) local dilation. On the other hand, (standard) local dilation always implies complex local dilation by letting $|j_1\rangle = 0$. We also note that, while it is always possible to take a basis in which the canonical

state is real (thanks to the Schmidt decomposition [NC10]), the above definition does not rely on the assumption of a real matrix representation of the canonical state.

In an alternative definition, a complex local dilation can be also understood as a *convex combination* of \tilde{S} and its complex conjugate, as introduced in [MNP21]. If we see local systems as a direct sum of subsystems:

$$\mathcal{H}_A = \mathcal{H}_{A_0} \oplus \mathcal{H}_{A_1}, \quad \mathcal{H}_B = \mathcal{H}_{B_0} \oplus \mathcal{H}_{B_1},$$

then the whole system satisfies

$$\mathcal{H}_{AB} = \mathcal{H}_A \otimes \mathcal{H}_B \cong \mathcal{H}_{A_0} \otimes \mathcal{H}_{B_0} \oplus \mathcal{H}_{A_0} \otimes \mathcal{H}_{B_1} \oplus \mathcal{H}_{A_1} \otimes \mathcal{H}_{B_0} \oplus \mathcal{H}_{A_1} \otimes \mathcal{H}_{B_1}.$$

If we only focus on vectors in the subspace

$$\mathcal{H}_{A_0} \otimes \mathcal{H}_{B_0} \oplus \mathcal{H}_{A_1} \otimes \mathcal{H}_{B_1} \subsetneq \mathcal{H}_A \otimes \mathcal{H}_B,$$

we then use the following diagonal direct sum notation for vectors $v_0 \in \mathcal{H}_{A_0} \otimes \mathcal{H}_{B_0}, v_1 \in \mathcal{H}_{A_1} \otimes \mathcal{H}_{B_1}$:

$$v_0 \oplus_{\Delta} v_1 := v_0 \oplus \vec{0}_{\mathcal{H}_{A_0} \otimes \mathcal{H}_{B_1}} \oplus \vec{0}_{\mathcal{H}_{A_1} \otimes \mathcal{H}_{B_0}} \oplus v_1 \in \mathcal{H}_{AB}.$$

That is, $v_0 \oplus_{\Delta} v_1$ is a vector in the subspace $\mathcal{H}_{A_0} \otimes \mathcal{H}_{B_0} \oplus \mathcal{H}_{A_1} \otimes \mathcal{H}_{B_1} \subsetneq \mathcal{H}_{AB}$.

Definition 2.5 (complex local dilation, alternative). *A strategy $\tilde{S} = (|\tilde{\psi}\rangle_{\tilde{A}\tilde{B}}, \{\tilde{E}_{xa}\}, \{\tilde{F}_{yb}\})$ is a complex local dilation of $S = (|\psi\rangle_{AB}, \{E_{xa}\}, \{F_{yb}\})$ if there exists local isometry $U = U_A \otimes U_B$ with*

$$\begin{aligned} U_A &: \mathcal{H}_A \rightarrow \mathcal{H}_{\tilde{A}_0} \otimes \mathcal{H}_{\tilde{A}'_0} \oplus \mathcal{H}_{\tilde{A}_1} \otimes \mathcal{H}_{\tilde{A}'_1}, \\ U_B &: \mathcal{H}_B \rightarrow \mathcal{H}_{\tilde{B}_0} \otimes \mathcal{H}_{\tilde{B}'_0} \oplus \mathcal{H}_{\tilde{B}_1} \otimes \mathcal{H}_{\tilde{B}'_1} \end{aligned}$$

such that

$$U[E_{xa} \otimes \text{Id}_B |\psi\rangle_{AB}] = (\tilde{E}_{xa} \otimes \text{Id}_{\tilde{B}_0} |\tilde{\psi}\rangle_{\tilde{A}_0\tilde{B}_0}) |j_0\rangle_{\hat{A}_0\hat{B}_0} \oplus_{\Delta} (\tilde{E}_{xa} \otimes \text{Id}_{\tilde{B}_1} |\tilde{\psi}\rangle_{\tilde{A}_1\tilde{B}_1}) |j_1\rangle_{\hat{A}_1\hat{B}_1}, \quad (4)$$

$$U[\text{Id}_A \otimes F_{yb} |\psi\rangle_{AB}] = (\text{Id}_{\tilde{A}_0} \otimes \tilde{F}_{yb} |\tilde{\psi}\rangle_{\tilde{A}_0\tilde{B}_0}) |j_0\rangle_{\hat{A}_0\hat{B}_0} \oplus_{\Delta} (\text{Id}_{\tilde{A}_1} \otimes \tilde{F}_{yb} |\tilde{\psi}\rangle_{\tilde{A}_1\tilde{B}_1}) |j_1\rangle_{\hat{A}_1\hat{B}_1} \quad (5)$$

hold for all a, b, x, y , where $|j_{0,1}\rangle$ are subnormalized state (not necessarily orthogonal): $\langle j_0|j_0\rangle + \langle j_1|j_1\rangle = 1$.

The following lemma shows the equivalence between the above definitions. In this paper we will primarily work with Definition 2.4.

Lemma 2.6. *Definitions 2.4 and 2.5 are equivalent.*

Proof. We show that Eq. (2) and Eq. (4) implies each other, and the rest can be proved similarly. In Eq. (4), we assume $\mathcal{H}_{\hat{A}_0\hat{B}_0} \cong \mathcal{H}_{\hat{A}_1\hat{B}_1}$ as we can always extend the smaller space to the larger one. Let

$$\begin{aligned} v_0 &:= (\tilde{E}_{xa} \otimes \text{Id}_{\tilde{B}_0} |\tilde{\psi}\rangle_{\tilde{A}_0\tilde{B}_0}) |j_0\rangle_{\hat{A}_0\hat{B}_0} \in \mathcal{H}_{\tilde{A}_0\tilde{B}_0\hat{A}_0\hat{B}_0} =: V_0, \\ v_1 &:= (\tilde{E}_{xa} \otimes \text{Id}_{\tilde{B}_1} |\tilde{\psi}\rangle_{\tilde{A}_1\tilde{B}_1}) |j_1\rangle_{\hat{A}_1\hat{B}_1} \in \mathcal{H}_{\tilde{A}_1\tilde{B}_1\hat{A}_1\hat{B}_1} =: V_1, \end{aligned}$$

then V_0 and V_1 has the same dimension, and $V_0 \oplus V_1 \cong V_0 \otimes \mathbb{C}^2$. Then

$$\text{Eq. (2)} = v_0 \oplus_{\Delta} v_1 \cong v_0 \otimes |0\rangle_{\mathbb{C}^2} + v_1 \otimes |1\rangle_{\mathbb{C}^2}, \text{Eq. (4)} = v_0 \otimes |00\rangle_{A'B'} + v_1 \otimes |11\rangle_{A'B'}.$$

By embedding \mathbb{C}^2 into $\mathcal{H}_{A'B'}$ as its subspace spanned by $\{|00\rangle, |11\rangle\}$ we have Eq. (2) \implies Eq. (4). Similarly, by projecting $\mathcal{H}_{A'B'}$ onto its subspace $\text{span}\{|00\rangle, |11\rangle\} \cong \mathbb{C}^2$, we have Eq. (4) \implies Eq. (2). \square

Now we can define *complex self-testing* in terms of complex local dilations.

Definition 2.7 (complex self-testing). *A strategy $\tilde{S} = (|\tilde{\psi}\rangle, \{\tilde{E}_{xa}\}, \{\tilde{F}_{yb}\})$ is complex self-tested by a correlation $p(a, b|x, y)$ if it is a complex local dilation of any strategy producing $p(a, b|x, y)$.*

Like its standard counterpart [BCK⁺25], one can introduce restrictions for the physical strategy and consider complex self-testing with assumptions. Specifically, we define

Definition 2.8 (complex self-testing with assumptions). *A strategy $\tilde{S} = (|\tilde{\psi}\rangle, \{\tilde{E}_{xa}\}, \{\tilde{F}_{yb}\})$ is complex full-rank self-tested by a correlation $p(a, b|x, y)$ if it is a complex local dilation of any full-rank strategy producing $p(a, b|x, y)$.*

A strategy $\tilde{S} = (|\tilde{\psi}\rangle, \{\tilde{E}_{xa}\}, \{\tilde{F}_{yb}\})$ is complex PVM self-tested by a correlation $p(a, b|x, y)$ if it is a complex local dilation of any projective strategy producing $p(a, b|x, y)$.

3 Basic properties, and lifting assumptions in complex self-testing

In this section we establish basic properties of complex local dilations that will be used in the remainder of this paper. Using them we will show that complex self-testing is free from PVM (the physical strategy performs projective measurement) and full-rank (the physical strategy employs a state with full Schmidt rank) assumptions, just like standard self-testing.²

The following two propositions show that a local dilation preserves (exact) support-preservingness and projectiveness.

Proposition 3.1 (Counterpart of Proposition 3.4 in [BCK⁺25]). *If $S \hookrightarrow_{\mathbb{C}} \tilde{S}$, then S is support preserving if and only if \tilde{S} is support-preserving.*

Proof. We will show support-preservingness via Lemma 2.2. In the remainder of this proof, we will construct operators \hat{E}_{xa} for each direction, and the construction of \hat{F}_{yb} is analogous.

The ‘If’ direction: Since \tilde{S} is support-preserving, there exist operators \hat{E}_{xa} such that $\tilde{E}_{xa} \otimes \text{Id} |\tilde{\psi}\rangle = \text{Id} \otimes \hat{E}_{xa} |\tilde{\psi}\rangle$ for all x, a , and therefore $\overline{\tilde{E}_{xa}} \otimes \text{Id} |\tilde{\psi}\rangle = \text{Id} \otimes \hat{E}_{xa} |\tilde{\psi}\rangle$. Construct operators

$$\hat{E}_{xa} := U_B^* [(\hat{E}_{xa} \otimes |0\rangle\langle 0|_{B'} + \overline{\hat{E}_{xa}} \otimes |1\rangle\langle 1|_{B'}) \otimes \text{Id}_{\hat{B}}] U_B, \quad \forall x, a.$$

²For a detailed discussion about assumptions in self-testing, see [BCK⁺25, Definition 2.3, 2.8]

Then,

$$\begin{aligned}
(\text{Id}_A \otimes \hat{E}_{xa}) |\psi\rangle &= (U_A^* U_A \otimes \hat{E}_{xa}) |\psi\rangle \\
&= (U_A^* \otimes U_B^*) (\text{Id}_{\hat{A}} \otimes \hat{E}_{xa}) |\tilde{\psi}\rangle |j_0\rangle |00\rangle + (\text{Id}_{\hat{A}} \otimes \overline{\hat{E}_{xa}}) |\tilde{\psi}\rangle |j_1\rangle |11\rangle \\
&= (U_A^* \otimes U_B^*) (U_A \otimes U_B) (E_{xa} \otimes \text{Id}_B) |\psi\rangle = (E_{xa} \otimes \text{Id}_B) |\psi\rangle.
\end{aligned}$$

Hence, S is support-preserving.

The ‘Only if’ direction: Since S is support-preserving, there exist operators \hat{E}_{xa} such that $E_{xa} \otimes \text{Id} |\psi\rangle = \text{Id} \otimes \hat{E}_{xa} |\psi\rangle$ for all x, a . Consider operators

$$\hat{E}_{xa} := U_B \hat{E}_{xa} U_B^* (|0\rangle\langle 0|_{B'} \otimes \text{Id}_{\hat{B}, \hat{B}}) + \overline{U_B \hat{E}_{xa} U_B^*} (|1\rangle\langle 1|_{B'} \otimes \text{Id}_{\hat{B}, \hat{B}}) \in L(\mathcal{H}_{\hat{B}, \hat{B}, B'}), \quad \forall x, a.$$

It holds that

$$\begin{aligned}
& (\text{Id}_{\hat{A}, \hat{A}, A'} \otimes \hat{E}_{xa}) |\tilde{\psi}\rangle |j_0\rangle |00\rangle + (\text{Id}_{\hat{A}, \hat{A}, A'} \otimes \overline{\hat{E}_{xa}}) |\tilde{\psi}\rangle |j_1\rangle |11\rangle \\
&= (\text{Id}_{\hat{A}, \hat{A}, A'} \otimes U_B \hat{E}_{xa} U_B^*) (|\tilde{\psi}\rangle |j_0\rangle |00\rangle + |\tilde{\psi}\rangle |j_1\rangle |11\rangle) \\
&= (\text{Id}_{\hat{A}, \hat{A}, A'} \otimes U_B \hat{E}_{xa} U_B^*) U |\psi\rangle \\
&= U (\text{Id}_A \otimes \hat{E}_{xa}) |\psi\rangle \\
&= U (E_{xa} \otimes \text{Id}_B) |\psi\rangle \\
&= (\tilde{E}_{xa} \otimes \text{Id}_{\hat{B}} |\tilde{\psi}\rangle) |j_0\rangle |00\rangle + (\overline{\tilde{E}_{xa}} \otimes \text{Id}_{\hat{B}} |\tilde{\psi}\rangle) |j_1\rangle |11\rangle.
\end{aligned}$$

So $(\text{Id}_{\hat{A}, \hat{A}, A'} \otimes \hat{E}_{xa}) |\tilde{\psi}\rangle |j_0\rangle |00\rangle = (\tilde{E}_{xa} \otimes \text{Id}_{\hat{B}} |\tilde{\psi}\rangle) |j_0\rangle |00\rangle$, which means that the operators $\tilde{E}_{xa} \otimes \text{Id}_{\hat{A}, A'}$ and $\Pi_{\hat{A}} \otimes \Pi_{j_0, A} \otimes |0\rangle\langle 0|_{A'}$ commute. Notice that

$$[\tilde{E}_{xa} \otimes \text{Id}_{\hat{A}, A'}, \Pi_{\hat{A}} \otimes \Pi_{j_0, A} \otimes |0\rangle\langle 0|_{A'}] = [\tilde{E}_{xa}, \Pi_{\hat{A}}] \otimes \Pi_{j_0, A} \otimes |0\rangle\langle 0|_{A'}.$$

Hence, \tilde{E}_{xa} and $\Pi_{\hat{A}}$ commute, so \tilde{S} is support-preserving. \square

Proposition 3.2 (Counterpart of Proposition 3.6 in [BCK⁺25]). *If $S \hookrightarrow_{\mathbb{C}} \tilde{S}$, then S is 0-projective if and only if \tilde{S} is 0-projective.*

Proof. Note that

$$\begin{aligned}
U[E_{xa} \otimes \text{Id}_B |\psi\rangle_{AB}] &= (\tilde{E}_{xa} \otimes \text{Id}_{\hat{B}} |\tilde{\psi}\rangle_{\hat{A}\hat{B}}) |j_0\rangle_{\hat{A}\hat{B}} |00\rangle_{A'B'} \\
&\quad + (\overline{\tilde{E}_{xa}} \otimes \text{Id}_{\hat{B}} |\tilde{\psi}\rangle_{\hat{A}\hat{B}}) |j_1\rangle_{\hat{A}\hat{B}} |11\rangle_{A'B'}, \\
U[(\text{Id}_A - E_{xa}) \otimes \text{Id}_B |\psi\rangle_{AB}] &= ((\text{Id}_{\hat{A}} - \tilde{E}_{xa}) \otimes \text{Id}_{\hat{B}} |\tilde{\psi}\rangle_{\hat{A}\hat{B}}) |j_0\rangle_{\hat{A}\hat{B}} |00\rangle_{A'B'} \\
&\quad + ((\text{Id}_{\hat{A}} - \overline{\tilde{E}_{xa}}) \otimes \text{Id}_{\hat{B}} |\tilde{\psi}\rangle_{\hat{A}\hat{B}}) |j_1\rangle_{\hat{A}\hat{B}} |11\rangle_{A'B'}.
\end{aligned}$$

Taking the inner product of the above two equations results in

$$\begin{aligned}
\langle \psi | E_{xa} (\text{Id}_A - E_{xa}) \otimes \text{Id}_B |\psi\rangle &= \langle \tilde{\psi} | \tilde{E}_{xa} (\text{Id}_{\hat{A}} - \tilde{E}_{xa}) \otimes \text{Id}_{\hat{B}} |\tilde{\psi}\rangle \langle j_0 | j_0 \rangle \\
&\quad + \langle \tilde{\psi} | \overline{\tilde{E}_{xa}} (\text{Id}_{\hat{A}} - \overline{\tilde{E}_{xa}}) \otimes \text{Id}_{\hat{B}} |\tilde{\psi}\rangle \langle j_1 | j_1 \rangle.
\end{aligned}$$

On the one hand, if $\langle \tilde{\psi} | \tilde{E}_{xa}(\text{Id}_{\tilde{A}} - \tilde{E}_{xa}) \otimes \text{Id}_{\tilde{B}} | \tilde{\psi} \rangle = 0$, then likewise $\langle \overline{\tilde{\psi}} | \overline{\tilde{E}_{xa}(\text{Id}_{\tilde{A}} - \tilde{E}_{xa})} \otimes \text{Id}_{\tilde{B}} | \overline{\tilde{\psi}} \rangle = 0$, and hence $\langle \psi | E_{xa}(\text{Id}_A - E_{xa}) \otimes \text{Id}_B | \psi \rangle = 0$. On the other hand, since both $\tilde{E}_{xa}(\text{Id}_{\tilde{A}} - \tilde{E}_{xa})$ and $(\text{Id}_{\tilde{A}} - \tilde{E}_{xa})$ are positive semidefinite, $\langle \psi | E_{xa}(\text{Id}_A - E_{xa}) \otimes \text{Id}_B | \psi \rangle = 0$ implies $\langle \tilde{\psi} | \tilde{E}_{xa}(\text{Id}_{\tilde{A}} - \tilde{E}_{xa}) \otimes \text{Id}_{\tilde{B}} | \tilde{\psi} \rangle = 0$, hence \tilde{S} is 0-projective. Thus, we conclude that S is 0-projective if and only if \tilde{S} is 0-projective. \square

With Propositions 3.2 and 3.1 at hand, we are ready to present the “lifting assumption” theorems.

Theorem 3.3 (Counterpart of Theorem 4.3 in [BCK⁺25]). *Let $p(a, b|x, y)$ be a quantum correlation. If $p(a, b|x, y)$ pure PVM complex self-tests a full-rank canonical strategy \tilde{S} , then \tilde{S} is projective, and $p(a, b|x, y)$ pure complex self-tests \tilde{S} .*

Proof. For any S (with a pure state) that generates $p(a, b|x, y)$, consider its Naimark dilation S_{Naimark} . Since $p(a, b|x, y)$ pure PVM self-tests \tilde{S} , it holds that $S_{\text{Naimark}} \hookrightarrow_{\mathbb{C}} \tilde{S}$. By Proposition 3.2, \tilde{S} is 0-projective, thus projective (since it is full-rank).

Again for any S (with a pure state) that generates $p(a, b|x, y)$, its Naimark dilation satisfies $S_{\text{Naimark}} \hookrightarrow_{\mathbb{C}} \tilde{S}$. Note that \tilde{S} is assumed to be full-rank (thus support-preserving); then, S_{Naimark} is support-preserving by Proposition 3.1. Consequently, S is support-preserving by [BCK⁺25, Theorem 3.18]. Then, $S \hookrightarrow S_{\text{Naimark}}$ by [BCK⁺25, Proposition 3.17]. By transitivity, $S \hookrightarrow_{\mathbb{C}} \tilde{S}$. So we conclude that $p(a, b|x, y)$ also pure self-tests \tilde{S} . \square

Theorem 3.4 (Counterpart of Theorem 4.5 in [BCK⁺25]). *Let $p(a, b|x, y)$ be a quantum correlation. If $p(a, b|x, y)$ pure full-rank complex self-tests a PVM canonical strategy \tilde{S} , then \tilde{S} is support-preserving, and $p(a, b|x, y)$ pure complex self-tests \tilde{S} .*

Proof. For any S (with a pure state) that generates $p(a, b|x, y)$, consider its restriction to the support S_{res} . Since $p(a, b|x, y)$ pure full-rank self-tests \tilde{S} , it holds that $S_{\text{res}} \hookrightarrow_{\mathbb{C}} \tilde{S}$. By Proposition 3.1, \tilde{S} is support-preserving.

Again for any S (with a pure state) that generates $p(a, b|x, y)$, its restriction satisfies $S_{\text{Naimark}} \hookrightarrow_{\mathbb{C}} \tilde{S}$. Note that \tilde{S} is assumed to be projective (thus 0-projective); then, S_{res} is 0-projective by Proposition 3.2. Consequently, S is 0-projective by [BCK⁺25, Theorem 3.9]. Then, $S \hookrightarrow S_{\text{res}}$ by [BCK⁺25, Proposition 3.8]. By transitivity, $S \hookrightarrow_{\mathbb{C}} \tilde{S}$. So we conclude that $p(a, b|x, y)$ also pure self-tests \tilde{S} . \square

We remark that this paper considers strategies with a pure state, therefore lifting the purity assumption in complex self-testing is beyond its scope, which we leave for future work.

Finally, we discuss the real simulation of a quantum strategy [MMG09], a special type of complex local dilation that may be of independent interest. It shows that any quantum correlation admits a real quantum realization (all matrix entries are real) via complex local dilation. Let $|\pm i\rangle := (|0\rangle \pm i|1\rangle)/\sqrt{2}$ be the eigenstate of Pauli matrix σ_Y .

Definition 3.5 (real simulation). Let $S = (|\psi\rangle, \{E_{xa}\}, \{F_{yb}\})$ be a complex strategy. The real simulation S_R of S is defined as $S_R := (|\psi_R\rangle, \{E_{R,xa}\}, \{F_{R,yb}\})$, where

$$\begin{aligned} |\psi_R\rangle &:= (|\psi\rangle | +i + i\rangle + |\bar{\psi}\rangle | -i - i\rangle) / \sqrt{2}, \\ E_{R,xa} &:= E_{xa} \otimes | +i\rangle\langle +i| + \overline{E_{xa}} \otimes | -i\rangle\langle -i|, \\ F_{R,yb} &:= F_{yb} \otimes | +i\rangle\langle +i| + \overline{F_{yb}} \otimes | -i\rangle\langle -i|. \end{aligned}$$

It is straightforward to verify that $|\psi_R\rangle, A_{R,xa}, B_{R,yb}$ all have real entries, and S_R gives the same correlation as S . By the definition of complex dilation 2.7 it holds that $S_R \hookrightarrow_{\mathbb{C}} S$. We remark that the auxiliary state $|\pm i\rangle$ is not strictly necessary; any state $|\phi\rangle$ satisfying $\langle \phi | \bar{\phi} \rangle = 0$ would suffice.

The following property relates real simulations of two strategies when one is a complex local dilations of another.

Proposition 3.6. *If $S \hookrightarrow_{\mathbb{C}} \tilde{S}$, then $S_R \hookrightarrow \tilde{S}_R$.*

Proof. Given that $S \hookrightarrow_{\mathbb{C}} \tilde{S}$, there exist local isometries V_A, V_B and auxiliary states $|j_0\rangle, |j_1\rangle$ that satisfy the complex local dilation relations. Now consider the action of $V_{A,R} := | +i\rangle\langle +i| \otimes V_A + | -i\rangle\langle -i| \otimes \overline{V_A}, V_{B,R} := V_B \otimes | +i\rangle\langle +i| + \overline{V_B} \otimes | -i\rangle\langle -i|$ on S_R . We have

$$\begin{aligned} & (V_{A,R} \otimes V_{B,R}) (E_{xa,R} \otimes F_{yb,R} |\psi_R\rangle) \\ &= (V_{A,R} \otimes V_{B,R}) \frac{1}{\sqrt{2}} (| +i + i\rangle (E_{xa} \otimes F_{yb} |\psi\rangle) + | -i - i\rangle (\overline{E_{xa}} \otimes \overline{F_{yb}} |\bar{\psi}\rangle)) \\ &= \frac{1}{\sqrt{2}} (| +i + i\rangle (V_A \otimes V_B) (E_{xa} \otimes F_{yb} |\psi\rangle) + | -i - i\rangle (\overline{V_A} \otimes \overline{V_B}) (\overline{E_{xa}} \otimes \overline{F_{yb}} |\bar{\psi}\rangle)) \\ &= \frac{1}{\sqrt{2}} (| +i + i\rangle |00\rangle |j_0\rangle (\tilde{E}_{xa} \otimes \tilde{F}_{yb} |\tilde{\psi}\rangle) + | +i + i\rangle |11\rangle |j_1\rangle (\tilde{\overline{E_{xa}}} \otimes \tilde{\overline{F_{yb}}} |\tilde{\bar{\psi}}\rangle) + \\ & \quad | -i - i\rangle |00\rangle |j_0\rangle (\tilde{\overline{E_{xa}}} \otimes \tilde{\overline{F_{yb}}} |\tilde{\bar{\psi}}\rangle) + | -i - i\rangle |11\rangle |j_1\rangle (\tilde{E}_{xa} \otimes \tilde{F}_{yb} |\tilde{\psi}\rangle)) \\ &= \frac{1}{\sqrt{2}} (|00\rangle |j_0\rangle) (| +i + i\rangle \tilde{E}_{xa} \otimes \tilde{F}_{yb} |\tilde{\psi}\rangle + | -i - i\rangle \tilde{\overline{E_{xa}}} \otimes \tilde{\overline{F_{yb}}} |\tilde{\bar{\psi}}\rangle) + \\ & \quad \frac{1}{\sqrt{2}} (|11\rangle |j_1\rangle) (| -i - i\rangle \tilde{E}_{xa} \otimes \tilde{F}_{yb} |\tilde{\psi}\rangle + | +i + i\rangle \tilde{\overline{E_{xa}}} \otimes \tilde{\overline{F_{yb}}} |\tilde{\bar{\psi}}\rangle). \end{aligned}$$

Let U_i be the 2-dimensional unitary that maps $|\pm i\rangle$ to $|\mp i\rangle$. Then consider the action of the local unitary $U := |0\rangle\langle 0| \otimes \text{Id}_j \otimes \text{Id}_i + |1\rangle\langle 1| \otimes \text{Id}_j \otimes U_i$. Clearly, $U \otimes U$ keeps $|00\rangle |j_0\rangle |\pm i \pm i\rangle$ unchanged, and maps $|11\rangle |j_1\rangle |\pm i \pm i\rangle$ to $|11\rangle |j_1\rangle |\mp i \mp i\rangle$. Therefore,

$$\begin{aligned} & (U \otimes \text{Id}_{\tilde{A}} \otimes U \otimes \text{Id}_{\tilde{B}}) (V_{A,R} \otimes V_{B,R}) (E_{xa,R} \otimes F_{yb,R} |\psi_R\rangle) \\ &= \frac{1}{\sqrt{2}} (|00\rangle |j_0\rangle) (| +i + i\rangle \tilde{E}_{xa} \otimes \tilde{F}_{yb} |\tilde{\psi}\rangle + | -i - i\rangle \tilde{\overline{E_{xa}}} \otimes \tilde{\overline{F_{yb}}} |\tilde{\bar{\psi}}\rangle) + \\ & \quad \frac{1}{\sqrt{2}} (|11\rangle |j_1\rangle) (| +i + i\rangle \tilde{E}_{xa} \otimes \tilde{F}_{yb} |\tilde{\psi}\rangle + | -i - i\rangle \tilde{\overline{E_{xa}}} \otimes \tilde{\overline{F_{yb}}} |\tilde{\bar{\psi}}\rangle) \\ &= (|00\rangle |j_0\rangle + |11\rangle |j_1\rangle) (\tilde{E}_{xa,R} \otimes \tilde{F}_{yb,R} |\tilde{\psi}_R\rangle). \end{aligned}$$

Also notice that $(|00\rangle |j_0\rangle + |11\rangle |j_1\rangle)$ is a unit vector. We conclude that $S_R \hookrightarrow \tilde{S}_R$ via local isometry $(U \otimes \text{Id}_{\tilde{A}} \otimes U \otimes \text{Id}_{\tilde{B}}) (V_{A,R} \otimes V_{B,R})$. \square

Finally, we point out that $S \hookrightarrow S_R$ fails in general for a full-rank S . This can be seen from the fact that strategies connected by a local dilation have the same higher moments ([PSZZ24, Proposition 4.8]), which are necessarily real-valued for S_R . This is possible only if S has real-valued moments.

4 An operator-algebraic characterization

Here we present the operator-algebraic picture of complex self-testing. We first show that a complex self-test is equivalent to the reference strategy having a unique real part of higher moments. Then, we describe this in terms of real states on real C* algebras.

4.1 Complex self-testing implies a unique real part of moments

We start with the simpler direction of the aforementioned equivalence.

Proposition 4.1. *If a support-preserving strategy \tilde{S} is complex self-tested by $p(a, b|x, y)$, then for any $k, \ell \in \mathbb{Z}^+$, for any \vec{x}, \vec{a} of length k , and any \vec{y}, \vec{b} of length ℓ ,*

$$\text{re} \langle \psi | E_{\vec{x}\vec{a}} \otimes F_{\vec{y}\vec{b}} | \psi \rangle$$

is the same across all strategies S producing $p(a, b|x, y)$.

Proof. Let $U = U_A \otimes U_B$ be the isometry and $|j_{0,1}\rangle$ be the auxiliary state from the complex self-test. We will prove that

$$U(E_{\vec{x}\vec{a}} \otimes \text{Id}_B) |\psi\rangle = (\tilde{E}_{\vec{x}\vec{a}} \otimes \text{Id}_{\tilde{B}} |\tilde{\psi}\rangle) |j_0\rangle |00\rangle + (\overline{\tilde{E}_{\vec{x}\vec{a}}} \otimes \text{Id}_{\tilde{B}} |\tilde{\psi}\rangle) |j_1\rangle |11\rangle \quad (6)$$

holds for all \vec{x}, \vec{a} by induction. First, from complex self-testing, for any single x, a it holds that

$$\begin{aligned} & (\tilde{E}_{xa} \otimes \text{Id}_{\tilde{B}} |\tilde{\psi}\rangle_{\tilde{A}\tilde{B}}) |j_0\rangle_{\hat{A}\hat{B}} |00\rangle_{A'B'} + (\overline{\tilde{E}_{xa}} \otimes \text{Id}_{\tilde{B}} |\tilde{\psi}\rangle_{\tilde{A}\tilde{B}}) |j_1\rangle_{\hat{A}\hat{B}} |11\rangle_{A'B'} \\ &= (U_A \otimes U_B)(E_{xa} \otimes \text{Id}) |\psi\rangle \\ &= (U_A E_{xa} U_A^* \otimes \text{Id})(U_A \otimes U_B) |\psi\rangle \\ &= (U_A E_{xa} U_A^* \otimes \text{Id}) |\tilde{\psi}\rangle_{\tilde{A}\tilde{B}} |j_0\rangle_{\hat{A}\hat{B}} |00\rangle_{A'B'} + (U_A E_{xa} U_A^* \otimes \text{Id}) |\tilde{\psi}\rangle_{\tilde{A}\tilde{B}} |j_1\rangle_{\hat{A}\hat{B}} |11\rangle_{A'B'}. \end{aligned} \quad (7)$$

It is clear that Eq. (6) is true with words of length 0 or 1. Suppose it is true for words of length k . Then for any \vec{x}, \vec{a} of length k , we have that

$$\begin{aligned} & U(E_{x_{k+1}a_{k+1}} E_{\vec{x}\vec{a}} \otimes \text{Id}_B) |\psi\rangle \\ &= (U_A E_{x_{k+1}a_{k+1}} U_A^* \otimes \text{Id}) U(E_{\vec{x}\vec{a}} \otimes \text{Id}_B) |\psi\rangle \\ &= (U_A E_{x_{k+1}a_{k+1}} U_A^* \otimes \text{Id}) ((\tilde{E}_{\vec{x}\vec{a}} \otimes \text{Id}_{\tilde{B}} |\tilde{\psi}\rangle) |j_0\rangle |00\rangle + (\overline{\tilde{E}_{\vec{x}\vec{a}}} \otimes \text{Id}_{\tilde{B}} |\tilde{\psi}\rangle) |j_1\rangle |11\rangle) \\ &= (U_A E_{x_{k+1}a_{k+1}} U_A^* \otimes \text{Id}) ((\text{Id}_{\tilde{A}} \otimes \hat{E}_{\vec{x}\vec{a}} |\tilde{\psi}\rangle) |j_0\rangle |00\rangle + (\text{Id}_{\tilde{A}} \otimes \overline{\hat{E}_{\vec{x}\vec{a}}} |\tilde{\psi}\rangle) |j_1\rangle |11\rangle) \\ &= (\text{Id}_{\tilde{A}\tilde{A}'} \otimes \hat{E}_{\vec{x}\vec{a}} \otimes \text{Id}_{\tilde{B}\tilde{B}'}) (U_A E_{x_{k+1}a_{k+1}} U_A^* \otimes \text{Id}) (|\tilde{\psi}\rangle |j_0\rangle |00\rangle) \\ &+ (\text{Id}_{\tilde{A}\tilde{A}'} \otimes \overline{\hat{E}_{\vec{x}\vec{a}}} \otimes \text{Id}_{\tilde{B}\tilde{B}'}) (U_A E_{x_{k+1}a_{k+1}} U_A^* \otimes \text{Id}) (|\tilde{\psi}\rangle |j_1\rangle |11\rangle). \end{aligned}$$

The third equation uses the fact that \tilde{S} is support preserving and Lemma 2.2. Multiplying both sides of Eq. (7) by $(\hat{E}_{\tilde{x}\tilde{a}} \otimes |0\rangle\langle 0|_{B'} + \overline{\hat{E}_{\tilde{x}\tilde{a}}} \otimes |1\rangle\langle 1|_{B'}) \otimes \text{Id}_{\tilde{A}\tilde{A}'\tilde{B}}$, we get

$$\begin{aligned}
& (\text{Id}_{\tilde{A}\tilde{A}'} \otimes \hat{E}_{\tilde{x}\tilde{a}} \otimes \text{Id}_{\tilde{B}B'}) (U_A E_{x_{k+1}a_{k+1}} U_A^* \otimes \text{Id}) (|\tilde{\psi}\rangle |j_0\rangle |00\rangle) \\
& + (\text{Id}_{\tilde{A}\tilde{A}'} \otimes \overline{\hat{E}_{\tilde{x}\tilde{a}}} \otimes \text{Id}_{\tilde{B}B'}) (U_A E_{x_{k+1}a_{k+1}} U_A^* \otimes \text{Id}) (|\tilde{\psi}\rangle |j_1\rangle |11\rangle) \\
& = (\tilde{E}_{x_{k+1}a_{k+1}} \otimes \hat{E}_{\tilde{x}\tilde{a}} |\tilde{\psi}\rangle) |j_0\rangle |00\rangle + (\overline{\tilde{E}_{x_{k+1}a_{k+1}}} \otimes \overline{\hat{E}_{\tilde{x}\tilde{a}}} |\tilde{\psi}\rangle) |j_1\rangle |11\rangle \\
& = (\tilde{E}_{x_{k+1}a_{k+1}} \tilde{E}_{\tilde{x}\tilde{a}} \otimes \text{Id}_{\tilde{B}} |\tilde{\psi}\rangle) |j_0\rangle |00\rangle + (\overline{\tilde{E}_{x_{k+1}a_{k+1}}} \overline{\tilde{E}_{\tilde{x}\tilde{a}}} \otimes \text{Id}_{\tilde{B}} |\tilde{\psi}\rangle) |j_1\rangle |11\rangle
\end{aligned}$$

Hence, Eq. (6) holds for all words. We can prove similar statement for Bob's operator. Then,

$$\begin{aligned}
U(E_{\tilde{x}\tilde{a}} \otimes F_{\tilde{y}\tilde{b}}) |\psi\rangle & = (U_A E_{\tilde{x}\tilde{a}} U_A^* \otimes \text{Id}_{\tilde{B}\tilde{B}B'}) U(\text{Id} \otimes F_{\tilde{y}\tilde{b}}) |\psi\rangle \\
& = (U_A E_{\tilde{x}\tilde{a}} U_A^* \otimes \text{Id}_{\tilde{B}\tilde{B}B'}) [(\text{Id}_{\tilde{A}} \otimes \tilde{F}_{\tilde{y}\tilde{b}}) |\tilde{\psi}\rangle |j_0\rangle |00\rangle + (\text{Id}_{\tilde{A}} \otimes \overline{\tilde{F}_{\tilde{y}\tilde{b}}}) |\tilde{\psi}\rangle |j_1\rangle |11\rangle] \\
& = (\text{Id}_{\tilde{A}\tilde{A}'} \otimes \tilde{F}_{\tilde{y}\tilde{b}}) (U_A E_{\tilde{x}\tilde{a}} U_A^* \otimes \text{Id}_{\tilde{B}\tilde{B}B'}) |\tilde{\psi}\rangle |j_0\rangle |00\rangle \\
& + (\text{Id}_{\tilde{A}\tilde{A}'} \otimes \overline{\tilde{F}_{\tilde{y}\tilde{b}}}) (U_A E_{\tilde{x}\tilde{a}} U_A^* \otimes \text{Id}_{\tilde{B}\tilde{B}B'}) |\tilde{\psi}\rangle |j_1\rangle |11\rangle \\
& = (\tilde{E}_{\tilde{x}\tilde{a}} \otimes \tilde{F}_{\tilde{y}\tilde{b}}) |\tilde{\psi}\rangle |j_0\rangle |00\rangle + (\overline{\tilde{E}_{\tilde{x}\tilde{a}}} \otimes \overline{\tilde{F}_{\tilde{y}\tilde{b}}}) |\tilde{\psi}\rangle |j_1\rangle |11\rangle.
\end{aligned}$$

Note that $U |\psi\rangle = |\tilde{\psi}\rangle |j_0\rangle |00\rangle + |\overline{\tilde{\psi}}\rangle |j_1\rangle |11\rangle$. Taking the inner product of the two sides respectively, we get

$$\langle \psi | E_{\tilde{x}\tilde{a}} \otimes F_{\tilde{y}\tilde{b}} | \psi \rangle = \langle \tilde{\psi} | \tilde{E}_{\tilde{x}\tilde{a}} \otimes \tilde{F}_{\tilde{y}\tilde{b}} | \tilde{\psi} \rangle |j_0\rangle|^2 + \langle \overline{\tilde{\psi}} | \overline{\tilde{E}_{\tilde{x}\tilde{a}}} \otimes \overline{\tilde{F}_{\tilde{y}\tilde{b}}} | \overline{\tilde{\psi}} \rangle |j_1\rangle|^2. \quad (8)$$

The final statement is achieved by taking the real part of both sides of the equation. \square

We remark that Eq. (8) also implies that if the canonical strategy \tilde{S} has all higher moments real, then so does S , in which case complex self-testing reduces to standard self-testing.

4.2 Unique real part of moments implies complex self-testing

Next, we prove the converse of Proposition 4.1, given an extreme $p(a, b|x, y)$.

Proposition 4.2. *Suppose a correlation $p(a, b|x, y)$ is extreme in C_q . If all strategies producing $p(a, b|x, y)$ have the same real parts of their moments, then there is a canonical \tilde{S} such that \tilde{S} is complex self-tested by correlation $p(a, b|x, y)$.*

The proof of Proposition 4.2 relies on the following two lemmas about real polynomials on $E_{xa} \otimes F_{yb}$ and irreducible strategies.

Lemma 4.3. *Let S, S' be irreducible strategies with the same real parts of their moments, and let f be any real polynomial. Then $f(E_{xa} \otimes F_{yb}) = 0$ if and only if $f(E'_{xa} \otimes F'_{yb}) = 0$.*

Consequently, if $g(E_{xa} \otimes F_{yb}) = \text{ild}$ for some real polynomial g , then $g(E'_{xa} \otimes F'_{yb}) = \pm \text{ild}$.

Proof. Due to the symmetry, it suffice to show that $f(E_{xa} \otimes F_{yb}) = 0$ implies $f(E'_{xa} \otimes F'_{yb}) = 0$. For any real self-adjoint polynomial g ,

$$\langle \psi | g(E_{xa} \otimes F_{yb}) | \psi \rangle = \text{re} \langle \psi | g(E_{xa} \otimes F_{yb}) | \psi \rangle = \text{re} \langle \psi' | g(E'_{xa} \otimes F'_{yb}) | \psi' \rangle = \langle \psi' | g(E'_{xa} \otimes F'_{yb}) | \psi' \rangle \quad (9)$$

by the assumption and self-adjointness. Let h be a real polynomial. Since $h^* f^* f h$ is a real self-adjoint polynomial, Eq. (9) and positive semidefiniteness of $(h^* f^* f h)(E'_{xa} \otimes F'_{yb})$ imply

$$\begin{aligned} \langle \psi | (h^* f^* f h)(E_{xa} \otimes F_{yb}) | \psi \rangle &= 0 \\ \implies \langle \psi' | (h^* f^* f h)(E'_{xa} \otimes F'_{yb}) | \psi' \rangle &= 0 \\ \implies (fh)(E'_{xa} \otimes F'_{yb}) | \psi' \rangle &= 0. \end{aligned}$$

Therefore, $f(E'_{xa} \otimes F'_{yb}) \cdot h(E'_{xa} \otimes F'_{yb}) | \psi' \rangle = 0$ holds for every real polynomial h . Since S' is irreducible, the set $\{h(E'_{xa} \otimes F'_{yb}) | \psi' \rangle : h \text{ a real polynomial}\}$ spans $\mathcal{H}_{A'} \otimes \mathcal{H}_{B'}$. Therefore, $f(E'_{xa} \otimes F'_{yb}) = 0$.

To prove the second part of the lemma, assume $g(E_{xa} \otimes F_{yb}) = \text{ild}$. Then,

$$(g^2 + 1)(E_{xa} \otimes F_{yb}) = 0, \quad (gh - hg)(E_{xa} \otimes F_{yb}) = 0 \quad \text{for all real polynomials } h.$$

Note that $g^2 + 1$ and $gh - hg$ are real polynomials. By the first part of the lemma,

$$(g^2 + 1)(E'_{xa} \otimes F'_{yb}) = 0, \quad (gh - hg)(E'_{xa} \otimes F'_{yb}) = 0 \quad \text{for all real polynomials } h.$$

Since S' is irreducible, only scalar operators commute with all $E'_{xa} \otimes F'_{yb}$, hence $g(E'_{xa} \otimes F'_{yb}) = \alpha \text{ild}$ for some $\alpha \in \mathbb{C}$. Finally, $(g^2 + 1)(E'_{xa} \otimes F'_{yb}) = 0$ implies $\alpha = \pm i$. \square

Lemma 4.4. *If two irreducible strategies S, S' have the same real parts of their moments, then they S' is unitarily equivalent to S or \bar{S} .*

Proof. Let $\mathcal{A} \subset L(\mathcal{H}_A) \otimes L(H_B)$ and $\mathcal{A}' \subset L(\mathcal{H}_{A'}) \otimes L(H_{B'})$ be unital real subalgebras generated by $E_{xa} \otimes F_{yb}$ and $E'_{xa} \otimes F'_{yb}$, respectively. By irreducibility,

$$L(\mathcal{H}_A) \otimes L(H_B) = \mathcal{A} + i\mathcal{A}, \quad L(\mathcal{H}_{A'}) \otimes L(H_{B'}) = \mathcal{A}' + i\mathcal{A}', \quad (10)$$

though these sums may not be direct sums. By Lemma 4.3, the map

$$\phi : \mathcal{A} \rightarrow \mathcal{A}', \quad f(E_{xa} \otimes F_{yb}) \mapsto f(E'_{xa} \otimes F'_{yb})$$

for real polynomials f is well-defined, and is an isomorphism of real $*$ -algebras. We distinguish two cases.

- Suppose that $f(E_{xa} \otimes F_{yb}) \neq \text{ild}$ for all real polynomials f . Then, both sums in (10) are direct sums by Lemma 4.3. Hence, we can extend ϕ to an isomorphism of complex $*$ -algebras $\phi : L(\mathcal{H}_A) \otimes L(H_B) \rightarrow L(\mathcal{H}_{A'}) \otimes L(H_{B'})$ via $\phi(X + iY) = \phi(X) + i\phi(Y)$ for $X, Y \in \mathcal{A}$.
- Suppose that $f(E_{xa} \otimes F_{yb}) = \text{ild}$ for some real polynomial f . Then, $f(E'_{xa} \otimes F'_{yb}) = \pm \text{ild}$ by Lemma 4.3. Consequently, $L(\mathcal{H}_A) \otimes L(H_B) = \mathcal{A}$ and $L(\mathcal{H}_{A'}) \otimes L(H_{B'}) = \mathcal{A}'$. Note that $\phi : L(\mathcal{H}_A) \otimes L(H_B)$ is an isomorphism of real $*$ -algebras. If $f(E'_{xa} \otimes F'_{yb}) = \text{ild}$, then $\phi(\text{ild}) = \text{ild}$, so ϕ is an isomorphism of complex $*$ -algebras; if $f(E'_{xa} \otimes F'_{yb}) = -\text{ild}$, then $\phi(\text{ild}) = -\text{ild}$, so $\bar{\phi}$ is an isomorphism of complex $*$ -algebras.

In both cases (after replacing S with \bar{S} in the second case if needed), there is an isomorphism of complex $*$ -algebras $\phi : L(\mathcal{H}_A) \otimes L(\mathcal{H}_B) \rightarrow L(\mathcal{H}_{A'}) \otimes L(\mathcal{H}_{B'})$ that maps $E_{xa} \otimes F_{yb}$ to $E'_{xa} \otimes F'_{yb}$. Note that ϕ maps $L(\mathcal{H}_A) \otimes \text{Id}$ onto $L(\mathcal{H}_{A'}) \otimes \text{Id}$ and $\text{Id} \otimes L(\mathcal{H}_B)$ onto $\text{Id} \otimes L(\mathcal{H}_{B'})$. Hence, by the Skolem-Noether theorem [Bre14, Theorem 4.46], there are unitaries $U : \mathcal{H}_A \rightarrow \mathcal{H}_{A'}$ and $V : \mathcal{H}_B \rightarrow \mathcal{H}_{B'}$ such that $\phi(Z) = (U \otimes V)Z(U^* \otimes V^*)$ for $Z \in L(\mathcal{H}_A \otimes \mathcal{H}_B)$. Finally, since

$$\langle \psi' | (U \otimes V)Z(U^* \otimes V^*) | \psi' \rangle = \langle \psi' | \phi(Z) | \psi' \rangle = \langle \psi | Z | \psi \rangle \quad \text{for all } Z \in L(\mathcal{H}_A \otimes \mathcal{H}_B),$$

it follows that $|\psi'\rangle = (U \otimes V)|\psi\rangle$. \square

Proof of Proposition 4.2. Given a quantum strategy $S = (|\psi\rangle, \{E_{xa}\}, \{F_{yb}\})$ producing $p(a, b|x, y)$, decompose it by the fundamental structure theorem of finite-dimensional C^* -algebras:

$$\begin{aligned} E_{xa} &= \bigoplus_i E_{xa}^{(i)} \otimes \text{Id} \in \bigoplus_i \mathcal{H}_A^{(i)} \otimes \mathcal{K}_A^{(i)}, \\ F_{yb} &= \bigoplus_j F_{yb}^{(j)} \otimes \text{Id} \in \bigoplus_j \mathcal{H}_B^{(j)} \otimes \mathcal{K}_B^{(j)}, \\ |\psi\rangle &= \bigoplus_{i,j} \left(\sum_{k,\ell} \sqrt{\lambda^{(ijk\ell)}} |\psi^{(ijk\ell)}\rangle \otimes |\alpha^{(ik)}, \beta^{(j\ell)}\rangle \right) \in \bigoplus_{i,j} \mathcal{H}_A^{(i)} \otimes \mathcal{H}_B^{(j)} \otimes \mathcal{K}_A^{(i)} \otimes \mathcal{K}_B^{(j)}, \end{aligned}$$

where positive coefficients $\lambda^{(ijk\ell)} > 0$ satisfies $\sum_{i,j,k,\ell} \lambda^{(ijk\ell)} = 1$, and $\{|\alpha^{(ik)}\rangle\}_k, \{|\beta^{(j\ell)}\rangle\}_\ell$ are orthonormal bases for $\mathcal{K}_A^{(i)}, \mathcal{K}_B^{(j)}$, respectively. Consider irreducible $S^{(ijk\ell)} := (|\psi^{(ijk\ell)}\rangle, \{E_{xa}^{(i)}\}, \{F_{yb}^{(j)}\})$ and their correlations $p^{(ijk\ell)}$. Then, $\sum_{i,j,k,\ell} \lambda^{(ijk\ell)} p^{(ijk\ell)} = p$. Since p is extreme in C_q , we have that $p^{(ijk\ell)} = p$. By our hypothesis, $S^{(ijk\ell)}$ have the same real parts of their moments. Thus, there exists an irreducible $\tilde{S} = (|\tilde{\psi}\rangle, \{\tilde{E}_{xa}\}, \{\tilde{F}_{yb}\})$ that produces p .

Since \tilde{S} and $S^{(ijk\ell)}$ agree on the real parts of their moments, $S^{(ijk\ell)}$ is unitarily equivalent to \tilde{S} or $\bar{\tilde{S}}$ by Lemma 4.4. As this holds for every quadruple $(ijk\ell)$, we conclude that $S \hookrightarrow_{\mathbb{C}} \tilde{S}$ by Definition 2.5. \square

4.3 The operator-algebraic formulation

In the language of operator algebras, having a unique real part of moments can be translated to uniqueness of a finite dimensional real state on a certain universal real C^* algebra. A real (unital) C^* algebra is a Banach $*$ -algebra over \mathbb{R} satisfying the C^* -identity $\|a^*a\| = \|a\|^2$, as well as the additional property that $1 + a^*a$ is invertible for every a (this ensures a real version of the GNS construction). The framework of real C^* algebras shares many similarities with that of complex C^* algebras, for instance the GNS construction. For a comprehensive introduction of real C^* algebras, see e.g., [Goo82, Li03]. A *universal* real C^* algebra $\mathcal{A}_{\mathbb{R}}(G, R)$ is a real C^* algebra such that (1) the elements of G generate $\mathcal{A}_{\mathbb{R}}(G, R)$ and satisfy the relations in R , and (2) has the following universal property: for any real C^* -algebra \mathcal{B} and any set of elements in \mathcal{B} that satisfy the same relations R , there exists a unique $*$ -homomorphism from $\mathcal{A}_{\mathbb{R}}(G, R)$ to \mathcal{B} that maps the generators accordingly.

Lemma 4.5. *The following statements are equivalent:*

1. For any $k, \ell \in \mathbb{N}^+$, for any \vec{x}, \vec{a} of length k and \vec{y}, \vec{b} of length ℓ , the real parts of moments

$$\operatorname{re} \langle \psi | E_{\vec{x}\vec{a}} \otimes F_{\vec{y}\vec{b}} | \psi \rangle$$

coincide for all strategy producing $p(a, b|x, y)$,

2. There is a unique finite dimensional real state on $\mathcal{A}_{\mathbb{R}, \text{POVM}}^{\mathcal{I}_A, \mathcal{O}_A} \otimes_{\min} \mathcal{A}_{\mathbb{R}, \text{POVM}}^{\mathcal{I}_B, \mathcal{O}_B}$ that agrees with $p(a, b|x, y)$.

Here, $\mathcal{A}_{\mathbb{R}, \text{POVM}}^{\mathcal{I}_A, \mathcal{O}_A}$ is the universal real C^* algebra generated by positive contractions $\{e_{xa} : x \in \mathcal{I}_A, a \in \mathcal{O}_A\}$, subject to the relations $\sum_a e_{xa} = 1, \forall x \in \mathcal{I}_A$, and similarly $\mathcal{A}_{\mathbb{R}, \text{POVM}}^{\mathcal{I}_B, \mathcal{O}_B}$ is generated by $\{f_{yb} : y \in \mathcal{I}_B, b \in \mathcal{O}_B\}$. A real state f agrees with $p(a, b|x, y)$ whenever $f(e_{xa} \otimes f_{yb}) = p(a, b|x, y)$ holds for all a, b, x, y .

Proof. (1) \Rightarrow (2): For any finite dimensional real state f that agrees with p , its real GNS construction [Li03, Theorem 3.3.4] gives a representation on a finite dimensional real Hilbert space, whose matrix representation gives raise to a real strategy which is moment-real. By Proposition 4.1, those f then agrees with all the words of generators, so f is determined on the whole real C^* algebra from its real linearity.

(2) \Rightarrow (1): Suppose $S^{(0)}, S^{(1)}$ differ in their real parts of moments, define real states f_0, f_1 by setting $f_0(e_{\vec{x}\vec{a}} \otimes f_{\vec{y}\vec{b}}) = \operatorname{re} \langle \psi^{(0)} | E_{\vec{x}\vec{a}}^{(0)} \otimes F_{\vec{y}\vec{b}}^{(0)} | \psi^{(0)} \rangle$, $f_1(e_{\vec{x}\vec{a}} \otimes f_{\vec{y}\vec{b}}) = \operatorname{re} \langle \psi^{(1)} | E_{\vec{x}\vec{a}}^{(1)} \otimes F_{\vec{y}\vec{b}}^{(1)} | \psi^{(1)} \rangle$, and extending them by real linearity. Then f_0, f_1 are valid real states on $\mathcal{A}_{\mathbb{R}, \text{POVM}}^{\mathcal{I}_A, \mathcal{O}_A} \otimes_{\min} \mathcal{A}_{\mathbb{R}, \text{POVM}}^{\mathcal{I}_B, \mathcal{O}_B}$ but $f_0 \neq f_1$. \square

We are ready to present the main result of this section.

Theorem 4.6.

1. If a support-preserving \tilde{S} is complex self-tested by a correlation $p(a, b|x, y)$, then there is a unique finite-dimensional real state on $\mathcal{A}_{\mathbb{R}, \text{POVM}}^{\mathcal{I}_A, \mathcal{O}_A} \otimes_{\min} \mathcal{A}_{\mathbb{R}, \text{POVM}}^{\mathcal{I}_B, \mathcal{O}_B}$ that agrees with $p(a, b|x, y)$.
2. Suppose the correlation $p(a, b|x, y)$ is extreme in C_q . If there is a unique finite-dimensional real state on $\mathcal{A}_{\mathbb{R}, \text{POVM}}^{\mathcal{I}_A, \mathcal{O}_A} \otimes_{\min} \mathcal{A}_{\mathbb{R}, \text{POVM}}^{\mathcal{I}_B, \mathcal{O}_B}$ that agrees with $p(a, b|x, y)$, then there is a canonical \tilde{S} such that \tilde{S} is complex self-tested by correlation $p(a, b|x, y)$.

Proof. Combine Propositions 4.1, 4.2, and Lemma 4.5. \square

Remark 4.7. For comparison with Theorem 4.6 above, consider the results of [PSZZ24], with notations slightly modified in accordance with ours (e.g., "centrally supported" in [PSZZ24] means "support-preserving" here). Namely, [PSZZ24, Proposition 4.10 and Theorem 4.12] establish the following.

1. If a support-preserving \tilde{S} is self-tested by a correlation $p(a, b|x, y)$, then there is a unique finite-dimensional state on $\mathcal{A}_{\text{POVM}}^{\mathcal{I}_A, \mathcal{O}_A} \otimes_{\min} \mathcal{A}_{\text{POVM}}^{\mathcal{I}_B, \mathcal{O}_B}$ that agrees with $p(a, b|x, y)$.

2. Given an correlation $p(a, b|x, y)$ that is extreme in C_q . If there is a unique finite-dimensional state on $\mathcal{A}_{POVM}^{I_A, O_A} \otimes_{\min} \mathcal{A}_{POVM}^{I_B, O_B}$ that agrees with $p(a, b|x, y)$, then there is a canonical \tilde{S} such that \tilde{S} is self-tested by correlation $p(a, b|x, y)$.

A reader can now readily compare the concepts of "self-test in terms of complex C^* algebra" and "complex self-test in terms of real C^* algebra". Note that $\mathcal{A}_{POVM}^{I, O} = \mathbb{C} \otimes_{\mathbb{R}} \mathcal{A}_{\mathbb{R}, POVM}^{I, O}$; alternatively, $\mathcal{A}_{\mathbb{R}, POVM}^{I_A, O_A}$ is the norm-closed real $*$ -subalgebra of $\mathcal{A}_{POVM}^{I_A, O_A}$ generated by the canonical generators of $\mathcal{A}_{POVM}^{I_A, O_A}$. Finally, we remark that Point 2 of Theorem 4.6 relies on the assumption that the considered correlation is extreme, similarly to its counterpart in [PSZZ24, Theorem 4.12]. This reflects a limitation of the existing proof techniques, and whether this assumption can be relaxed remains an interesting open question.

5 Realness of quantum strategies

Section 4 indicates that the real parts of higher moments are essential in complex self-testing, and leads our attention to quantum strategies with real moments. An obvious candidate of that is the family of strategies with a real matrix representation. Then the natural question to ask is, are there any other strategies with real higher moments? If the answer is affirmative then it would be a more appropriate definition of 'non-complex' quantum strategies in the context of self-testing. Here we solve this problem by fully identifying the family of strategies with real higher moments, which we will call 'self-conjugate' strategies.

To investigate whether a quantum strategy has real moments or admits a real matrix representation, one needs to consider the real algebra generated by measurements in the strategy. Given $X_1, \dots, X_m \in M_d(\mathbb{C})$ let $\text{Alg}_{\mathbb{R}}(X_j: j)$ and $\text{Alg}_{\mathbb{C}}(X_j: j)$ denote the real unital $*$ -algebra and the complex unital $*$ -algebra, respectively, generated by X_1, \dots, X_m in $M_d(\mathbb{C})$ endowed with the conjugate transpose. The collection X_1, \dots, X_m is *irreducible* if $\text{Alg}_{\mathbb{C}}(X_j: j) = M_d(\mathbb{C})$. In this case, $\text{Alg}_{\mathbb{R}}(X_j: j)$ is isomorphic to $M_d(\mathbb{R})$, $M_d(\mathbb{C})$ or $M_{d/2}(\mathbb{H})$ as a consequence of Frobenius' theorem [Goo82, Bre14].

Let us also record basic properties of the standard matrix representation of quaternions. Throughout the rest paper denote

$$J = \begin{pmatrix} 0 & -1 \\ 1 & 0 \end{pmatrix} \in M_2(\mathbb{R}). \quad (11)$$

Let $\mathbb{H} = \mathbb{R} + \mathbb{R}i + \mathbb{R}j + \mathbb{R}k$ denote the quaternion algebra (see e.g. [Bre14, Section 1.1]), and consider $n \times n$ quaternion matrices $M_n(\mathbb{H}) = \mathbb{H} \otimes_{\mathbb{R}} M_n(\mathbb{R})$ as a real $*$ -algebra, whose involution is the tensor product of the transpose in $M_n(\mathbb{R})$ and the canonical (symplectic) involution in \mathbb{H} . There is standard $*$ -embedding

$$\Phi : \mathbb{H} \rightarrow M_2(\mathbb{C}), \quad \alpha_0 + \alpha_1 i + \alpha_2 j + \alpha_3 k \mapsto \begin{pmatrix} \alpha_0 + \alpha_1 i & \alpha_2 + \alpha_3 i \\ -\alpha_2 + \alpha_3 i & \alpha_0 - \alpha_1 i \end{pmatrix}.$$

Note that $\text{ran } \Phi$ is, as a real algebra, generated by i -multiples of the Pauli matrices $i\sigma_X, i\sigma_Y, i\sigma_Z$. A direct calculation shows that

$$\overline{\Phi(z)} = J\Phi(z)J^* \quad (12)$$

for all $z \in \mathbb{H}$. Furthermore, Φ extends to the $*$ -embedding of real algebras

$$\Phi_n = \Phi \otimes_{\mathbb{R}} \text{Id}_{M_n(\mathbb{R})} : M_n(\mathbb{H}) = \mathbb{H} \otimes_{\mathbb{R}} M_n(\mathbb{R}) \hookrightarrow M_2(\mathbb{C}) \otimes_{\mathbb{R}} M_n(\mathbb{R}) = M_{2n}(\mathbb{C}).$$

Then $\text{ran } \Phi_n$ generates $M_{2n}(\mathbb{C})$ as a complex algebra, and $\text{Tr } X \in \mathbb{R}$ for every $X \in \text{ran } \Phi_n$. By (12), we have $\overline{X} = (J \otimes \text{Id}_n)X(J \otimes \text{Id}_n)^*$ for all $X \in \text{ran } \Phi_n$. This is a distinguishing feature of quaternionic matrices (as opposed to real and complex matrices): namely, in their irreducible representation on a Hilbert space, entry-wise complex conjugation coincides with conjugation by a unitary.

Proposition 5.1. *For an irreducible collection $X_1, \dots, X_m \in M_d(\mathbb{C})$, consider the following statements:*

1. *there is $U \in U_d(\mathbb{C})$ such that $UX_jU^* \in M_d(\mathbb{R})$ for $j = 1, \dots, m$;*
2. *there is $U \in U_d(\mathbb{C})$ such that $UX_jU^* = \overline{X_j}$ for $j = 1, \dots, m$;*
3. *$\text{Tr } X \in \mathbb{R}$ for every product X of X_1, \dots, X_m ;*
4. *$\text{Alg}_{\mathbb{R}}(X_j : j) \cap \mathbb{C}I = \mathbb{R}I$;*
5. *$\text{Alg}_{\mathbb{R}}(X_j : j) \neq M_d(\mathbb{C})$.*

Then (1) \Rightarrow (2) \Leftrightarrow (3) \Leftrightarrow (4) \Leftrightarrow (5).

If d is odd, or $d = 2$ and X_j are hermitian, or $d \in \{4, 6\}$ and $m \leq 3$ and X_j are projections, then (1) \Leftrightarrow (2) \Leftrightarrow (3) \Leftrightarrow (4) \Leftrightarrow (5).

Proof. The implications (2) \Rightarrow (3) \Rightarrow (4) \Rightarrow (5) are straightforward. Also, $UXU^* \in M_d(\mathbb{R})$ for $U \in U_d(\mathbb{C})$ implies $\overline{X} = (\overline{U}^*U)X(\overline{U}^*U)^*$ and $\overline{U}^*U \in U_d(\mathbb{C})$, so (1) \Rightarrow (2) holds.

Now assume (5) holds. Since $\text{Alg}_{\mathbb{R}}(X_j : j)$ is closed under the conjugate transpose, it is a semisimple real algebra. Furthermore, it is a simple real algebra since X_j are irreducible. Therefore $\text{Alg}_{\mathbb{R}}(X_j : j)$ is isomorphic to one of the $M_n(\mathbb{R})$, $M_n(\mathbb{C})$, $M_n(\mathbb{H})$ for some $n \in \mathbb{N}$ by [Bre14, Corollary 2.69]. Moreover, since the involution on $\text{Alg}_{\mathbb{R}}(X_j : j)$ is a restriction of the conjugate transpose and is therefore positive, the $*$ -algebra $\text{Alg}_{\mathbb{R}}(X_j : j)$ is $*$ -isomorphic to one of the real $*$ -algebras $M_n(\mathbb{R})$, $M_n(\mathbb{C})$, $M_n(\mathbb{H})$ with their canonical standard involutions for some $n \in \mathbb{N}$ by [PS76, Theorem 1.2]. Note that there is a canonical surjective $*$ -homomorphism of complex $*$ -algebras

$$\mathbb{C} \otimes_{\mathbb{R}} \text{Alg}_{\mathbb{R}}(X_j : j) \rightarrow \text{Alg}_{\mathbb{C}}(X_j : j) = M_d(\mathbb{C}), \quad (13)$$

and

$$\mathbb{C} \otimes_{\mathbb{R}} M_n(\mathbb{R}) \cong M_n(\mathbb{C}), \quad \mathbb{C} \otimes_{\mathbb{R}} M_n(\mathbb{C}) \cong M_n(\mathbb{C}) \times M_n(\mathbb{C}), \quad \mathbb{C} \otimes_{\mathbb{R}} M_n(\mathbb{H}) \cong M_{2n}(\mathbb{C}).$$

Suppose $\text{Alg}_{\mathbb{R}}(X_j : j) \cong M_n(\mathbb{C})$. Surjectivity of the homomorphism (13) implies $n = d$, and therefore $\text{Alg}_{\mathbb{R}}(X_j : j) = M_d(\mathbb{C})$, which contradicts (5). Therefore $\text{Alg}_{\mathbb{R}}(X_j : j)$ is $*$ -isomorphic to either $M_n(\mathbb{R})$ or $M_n(\mathbb{H})$. Since $M_n(\mathbb{C})$ and $M_{2n}(\mathbb{C})$ are simple algebras, the surjective homomorphism (13) is also injective, and therefore either $\text{Alg}_{\mathbb{R}}(X_j : j) \cong M_d(\mathbb{R})$, or d is even and $\text{Alg}_{\mathbb{R}}(X_j : j) \cong M_{d/2}(\mathbb{H})$.

In the first case, there is a $*$ -isomorphism of real algebras $\Psi : \text{Alg}_{\mathbb{R}}(X_j : j) \rightarrow M_d(\mathbb{R})$. Then $\text{Id}_{\mathbb{C}} \otimes_{\mathbb{R}} \Psi : M_d(\mathbb{C}) \rightarrow M_d(\mathbb{C})$ is an automorphism. By the Skolem-Noether theorem [Bre14, Theorem 1.30] there exists $V \in \text{GL}_d(\mathbb{C})$ such that $VX_jV^{-1} = \Psi(X_j)$ for $j = 1, \dots, m$. Since Ψ is a $*$ -homomorphism,

$$V^{-*}X_jV^* = (VX_j^*V^*)^* = \Psi(X_j^*)^* = \Psi(X_j) = VX_jV^{-1}$$

and therefore $X_jV^*V = V^*VX_j$ for $j = 1, \dots, m$. Since X_j are irreducible, it follows that $V^*V = \alpha I$ for some nonzero scalar α . Clearly $\alpha > 0$. Then $U = \frac{1}{\sqrt{\alpha}}V \in U_d(\mathbb{C})$ satisfies (1).

Now consider the second case. Then d is even and there is a $*$ -isomorphism of real algebras $\Psi : \text{Alg}_{\mathbb{R}}(X_j : j) \rightarrow M_{d/2}(\mathbb{H})$. Then $\text{Id}_{\mathbb{C}} \otimes_{\mathbb{R}} \Psi : M_d(\mathbb{C}) \rightarrow M_d(\mathbb{C})$ is again an automorphism. By the Skolem-Noether theorem [Bre14, Theorem 1.30] there exists $V \in \text{GL}_d(\mathbb{C})$ such that $VX_jV^{-1} = \Phi_n(\Psi(X_j))$ for $j = 1, \dots, m$, where $\Phi_n : M_{d/2}(\mathbb{H}) \hookrightarrow M_d(\mathbb{C})$ is the $*$ -embedding from Section 6.2. Since $\Phi_n \circ \Psi$ is a $*$ -homomorphism,

$$V^{-*}X_jV^* = (VX_j^*V^*)^* = ((\Phi_n \circ \Psi)(X_j^*))^* = (\Phi_n \circ \Psi)(X_j) = VX_jV^{-1}$$

and therefore $X_jV^*V = V^*VX_j$ for $j = 1, \dots, m$. Since X_j are irreducible, it follows that $V^*V = \alpha I$ for some nonzero scalar α . Clearly $\alpha > 0$. Then $W = \frac{1}{\sqrt{\alpha}}V \in U_d(\mathbb{C})$ satisfies $WX_jW^* = \Phi_n(\Psi(X_j))$ for $j = 1, \dots, m$. By (12),

$$\overline{WX_jW^*} = \overline{\Phi_n(\Psi(X_j))} = (J \otimes I_{d/2})\Phi_n(\Psi(X_j))(J \otimes I_{d/2})^* = (J \otimes I_{d/2})WX_jW^*(J \otimes I_{d/2})^*$$

and therefore $\overline{X_j} = UX_jU^*$ for $U = W^\top(J \otimes I_{d/2})W \in U_d(\mathbb{C})$, so (2) holds.

Finally, notice that the first case $\text{Alg}_{\mathbb{R}}(X_j : j) \cong M_d(\mathbb{R})$ is the only possibility whenever d is odd, or if $d = 2$ and X_j are hermitian matrices (since \mathbb{H} is not generated by hermitian elements), or if $d \in \{4, 6\}$, $m \leq 3$ and X_j are projections by Proposition 6.6 (which we will formally introduce later). \square

Recall that a finite-dimensional strategy S is irreducible if the $\{E_{xa}\}_{x,a}$ generate $L(\mathcal{H}_A)$ and the $\{F_{yb}\}_{y,b}$ generate $L(\mathcal{H}_B)$ as complex algebras.

Definition 5.2. *The strategy S is:*

1. (Schmidt) real if some (Schmidt) matrix representation of S is real;
2. (Schmidt) self-conjugate if for some/all (Schmidt) basis there exist local unitaries U_A, U_B such that

$$U_A E_{xa} U_A^* = \overline{E_{xa}}, \quad U_B F_{yb} U_B^* = \overline{F_{yb}}, \quad U_A \otimes U_B |\psi\rangle = |\overline{\psi}\rangle$$

holds for all x, y, a, b .

3. moment-real if $\langle \psi | E \otimes F | \psi \rangle \in \mathbb{R}$ for all words $E_{\vec{x}\vec{a}}$ of POVM operators $\{E_{xa}\}$ and words $F_{\vec{y}\vec{b}}$ of POVM operators $\{F_{yb}\}$.

Theorem 5.3. *For an irreducible strategy S with local Hilbert spaces \mathcal{H}_A and \mathcal{H}_B , consider the following statements:*

1. S is Schmidt real;
2. S is real;
3. S is Schmidt self-conjugate;
4. S is self-conjugate;
5. S is moment-real.

Then $(1) \Leftrightarrow (2) \Rightarrow (3) \Leftrightarrow (4) \Leftrightarrow (5)$.

If on Alice's side and Bob's side, at least one of the conditions

- local dimension is 2 or odd,
- local dimension is 4 or 6, there are at most three inputs, and measurements are binary and projective,

is fulfilled, then $(1) \Leftrightarrow (2) \Leftrightarrow (3) \Leftrightarrow (4) \Leftrightarrow (5)$.

Proof. The equivalences $(1) \Leftrightarrow (2)$ and $(3) \Leftrightarrow (4)$ follow from the existence of the singular value decomposition for real matrices, and the implications $(2) \Rightarrow (4) \Rightarrow (5)$ are straightforward. Let $d_A = \dim \mathcal{H}_A$ and $d_B = \dim \mathcal{H}_B$.

For the sake of contradiction, suppose that S is moment real but not self-conjugate. Assume that the conditions in Definition 5.2 fail for $\{E_{xa}\}$. Then by Proposition 5.1 $\{E_{xa}\}$ generate $L(\mathcal{H}_A)$ as a real algebra. Let $|\psi\rangle = \sum_{i=1}^r |u_i\rangle |v_i\rangle$ for linearly independent $|u_i\rangle \in \mathcal{H}_A$ and linearly independent $|v_i\rangle \in \mathcal{H}_B$. Note that

$$\langle \psi | E \otimes F | \psi \rangle = \sum_{i,j=1}^r \langle u_i | E | u_j \rangle \cdot \langle v_i | F | v_j \rangle \quad (14)$$

for $E \in L(\mathcal{H}_A)$ and $F \in L(\mathcal{H}_B)$. Since $\{F_{yb}\}$ are irreducible, there exists a word F of $\{F_{yb} : y, b\}$ such that not all $\langle v_i | F | v_j \rangle$ are 0. In particular, $\langle v_{i_0} | F | v_{j_0} \rangle \neq 0$ for some i_0, j_0 . Since $\{E_{xa}\}$ generate $L(\mathcal{H}_A)$ as a real algebra, there is a real combination $E = \sum_k \alpha_k E_k$ of words E_k of E_{xa} such that

$$\langle u_i | E | u_j \rangle = \begin{cases} i \langle v_{i_0} | F | v_{j_0} \rangle & \text{if } i = i_0, j = j_0 \\ 0 & \text{otherwise.} \end{cases}$$

Therefore

$$\sum_k \alpha_k \langle \psi | E_k \otimes F | \psi \rangle = \langle \psi | E \otimes F | \psi \rangle \notin \mathbb{R}$$

by Eq. (14), and so $\langle \psi | E_k \otimes F | \psi \rangle \notin \mathbb{R}$ for some k , which contradicts S being moment real.

Therefore there are $U_A \in U_{d_A}(\mathbb{C})$ and $U_B \in U_{d_B}(\mathbb{C})$ such that

$$U_A E_{xa} U_A^* = \overline{E_{xa}}, \quad U_B F_{yb} U_B^* = \overline{F_{yb}}$$

for all x, y, a, b . Denote $|\psi'\rangle = U_A^* \otimes U_B^* |\bar{\psi}\rangle$. Clearly,

$$\langle \psi' | E \otimes F | \psi' \rangle = \langle \bar{\psi} | U(E \otimes F) U^* | \bar{\psi} \rangle = \overline{\langle \psi | E \otimes F | \psi \rangle} = \langle \psi | E \otimes F | \psi \rangle \quad (15)$$

for all words E of E_{xa} and words F of F_{yb} . Since both sides of Eq. (15) are complex linear in $E \otimes F$, and $\text{Alg}_{\mathbb{C}}(E_{xa}: x, a) \otimes \text{Alg}_{\mathbb{C}}(F_{yb}: y, b) = L(\mathcal{H}_A \otimes \mathcal{H}_B)$, it furthermore follows that $\text{Tr}(|\psi'\rangle\langle\psi'|T) = \langle\psi'|T|\psi'\rangle = \langle\psi|T|\psi\rangle = \text{Tr}(|\psi\rangle\langle\psi|T)$ for all $T \in L(\mathcal{H}_A \otimes \mathcal{H}_B)$. Then $|\psi\rangle\langle\psi| = |\psi'\rangle\langle\psi'|$, and so $|\psi'\rangle = \alpha|\psi\rangle$ for some phase $\alpha \in \mathbb{C}$ of modulus 1. Therefore, we have

$$\overline{E_{xa}} = (\alpha U_A) E_{xa} (\alpha U_A)^*, \quad \overline{F_{yb}} = U_B F_{yb} U_B^*, \quad |\overline{\psi}\rangle = U_A \otimes U_B |\psi'\rangle = (\alpha U_A) \otimes U_B |\psi\rangle$$

for unitaries αU_A and U_B . Thus, (5) \Leftrightarrow (4).

Finally, assume that (4) holds, and that at least one of the exceptional conditions is fulfilled on Alice's and on Bob's side. By Proposition 5.1 there in particular exist orthonormal bases \mathcal{B}_A and \mathcal{B}_B relative to which the measurements in S are given by real matrices $E_{xa} \in M_{d_A}(\mathbb{R})$ and $F_{yb} \in M_{d_B}(\mathbb{R})$. Since S is irreducible, E_{xa} and F_{yb} generate $M_{d_A}(\mathbb{R})$ and $M_{d_B}(\mathbb{R})$ as real algebras. Therefore by (5),

$$\langle\psi|A \otimes B|\psi\rangle \in \mathbb{R}$$

for all $A \in M_{d_A}(\mathbb{R})$ and $B \in M_{d_B}(\mathbb{R})$. Write $|\psi\rangle = \sum_{i,j} \alpha_{ij} |ij\rangle$ relative to bases \mathcal{B}_A and \mathcal{B}_B ; then

$$\alpha_{ij} \overline{\alpha_{k\ell}} = \langle\psi|(|k\rangle\langle i| \otimes |j\rangle\langle\ell|)|\psi\rangle \in \mathbb{R}$$

for all $i, k = 1, \dots, d_A$ and $j, \ell = 1, \dots, d_B$. Therefore arguments of $\{\alpha_{ij}\}$ coincide, so there is $\zeta \in \mathbb{C}$ of modulus 1 such that $\zeta\{\alpha_{ij}\} \in \mathbb{R}^{d_A \times d_B}$. Therefore S is a real strategy, with corresponding orthonormal bases $\zeta\mathcal{B}_A$ and \mathcal{B}_B , and hence (2) holds. \square

Let us point out that there exist strategies that are moment-real but not real.

Example 5.4. *Let $d \geq 4$ be even. By Proposition 6.6, there exists an irreducible collection of projections $P_1, \dots, P_4 \in M_d(\mathbb{C})$ such that $\text{Tr}(P) \in \mathbb{R}$ for every product P of P_1, \dots, P_4 , and there is no $U \in U_d(\mathbb{C})$ such that $UP_j U^* \in M_d(\mathbb{R})$ for all $j = 1, \dots, 4$. Namely, for P_j one can take any projective generators of $M_{d/2}(\mathbb{H})$ within $M_d(\mathbb{C})$ (if $d \geq 8$, three projections suffice). Let*

$$S = \left(|\phi_d\rangle, \{P_j, \text{Id} - P_j\}_{j=1}^4, \{P_j^\top, \text{Id} - P_j^\top\}_{j=1}^4 \right)$$

where $|\phi_d\rangle = \frac{1}{\sqrt{d}} \sum_{i=1}^d |ii\rangle$ is the canonical maximally entangled state. Then the strategy S is self-conjugate but not real.

As seen in the above arguments, the distinction between real and self-conjugate strategies essentially boils down to the fact that realness of trace cannot distinguish between real matrices and quaternion matrices. Nevertheless, tracial identities distinguish between $M_m(\mathbb{R})$ and $M_n(\mathbb{H})$ for all m, n by [Row80, Corollary 2.5.12 and Remark 2.5.1] and [KvV18, Proposition 2.3].

Given the result of Theorem 5.3, we shall call a strategy S *non-real* if S is not real (or not Schmidt real). We shall call S *complex* if S is not self-conjugate (or not moment-real). We will show that it is exactly the class of complex strategies that cannot be self-tested (in the standard sense).

Theorem 5.5. *Let $\tilde{S} = (|\tilde{\psi}\rangle, \{\tilde{E}_{xa}\}, \{\tilde{F}_{yb}\})$ be a full-rank complex PVM strategy. Then \tilde{S} is not self-tested.*

Proof. Not all moments of \tilde{S} are real, the strategy \tilde{S} and its complex conjugate give rise to distinct states on $\mathcal{A}_{\mathbb{C},\text{POVM}}^{\mathcal{I}_A, \mathcal{O}_A} \otimes_{\min} \mathcal{A}_{\mathbb{C},\text{POVM}}^{\mathcal{I}_B, \mathcal{O}_B}$. Thus, the correlation of \tilde{S} is not an abstract self-test, and thus not a self-test by [PSZZ24, Proposition 4.10]. \square

Lemma 5.6. *Let S be a complex strategy and S_R be its real simulation. Then $S \hookrightarrow S_R$ does not hold.*

Proof. If $S \hookrightarrow S_R$ then S and S_R have the same moments. But S_R is moment-real, while S is not. \square

In the end of this section, we point out that for a full-rank, PVM strategy to be complex self-tested, it cannot be merely ‘one-sided real’.

Theorem 5.7. *Let $\tilde{S} = (|\tilde{\psi}\rangle, \{\tilde{E}_{xa}\}, \{\tilde{F}_{yb}\})$ be a full-rank PVM strategy. If there exist a basis where*

- $|\tilde{\psi}\rangle$ has real matrix representation,
- $\{\tilde{E}_{xa}\}$ has real matrix representation,
- at least one of $\{\tilde{F}_{yb}\}$ has no real matrix representation,

then \tilde{S} is not complex self-tested.

Proof. We prove this by showing $\text{re } \tilde{S} := (|\tilde{\psi}\rangle, \{\tilde{E}_{xa}\}, \{\text{re } \tilde{F}_{yb}\})$ produces the same correlation as \tilde{S} , but cannot be complex local dilated to \tilde{S} . First, we note that $\text{re } \tilde{F}_{yb} = 1/2(\tilde{F}_{yb} + \overline{\tilde{F}_{yb}})$, a convex combination of POVMs. So $\{\text{re } \tilde{F}_{yb}\}_b$ is a valid POVM.

To show that $\text{re } \tilde{S}$ produces the same correlation as \tilde{S} , notice that $\text{im } \tilde{F}_{yb}$ is anti-symmetric. Therefore $\langle \tilde{\psi} | \tilde{E}_{xa} \otimes \text{im } \tilde{F}_{yb} | \tilde{\psi} \rangle = 0$ for all a, b, x, y . Then $\langle \tilde{\psi} | \tilde{E}_{xa} \otimes \tilde{F}_{yb} | \tilde{\psi} \rangle = \langle \tilde{\psi} | \tilde{E}_{xa} \otimes \text{re } \tilde{F}_{yb} | \tilde{\psi} \rangle$.

To show that $\text{re } \tilde{S}$ cannot be complex local dilated to \tilde{S} , we first prove that for any measurement F_{yb} , $\{\text{re } F_{yb}\}$ is a PVM if and only if $F_{yb} = \overline{F_{yb}}$. We have that

$$\text{re } F_{yb}^2 = \frac{1}{4}(F_{yb} + \overline{F_{yb}} + \overline{F_{yb}}F_{yb} + F_{yb}\overline{F_{yb}}).$$

So $\text{re } F_{yb}$ being projection $((\text{re } F_{yb})^2 = \text{re } F_{yb})$ is equivalent to

$$\begin{aligned} \overline{F_{yb}}F_{yb} + F_{yb}\overline{F_{yb}} &= \overline{F_{yb}} + F_{yb} \\ \Leftrightarrow F_{yb}\overline{F_{yb}}F_{yb} &= F_{yb}. \end{aligned}$$

Also note that both $F_{yb}, \overline{F_{yb}}$ are projections of the same rank. So this implies $F_{yb} = \overline{F_{yb}}$.

Hence given the assumption, $\text{re } \tilde{S}$ is not a PVM. Therefore $\text{re } \tilde{S} \hookrightarrow_{\mathbb{C}} \tilde{S}$ does not hold, because complex local dilation preserves projectivity (Proposition 3.2). \square

6 A quaternion middle ground possibility

Proposition 5.5 shows that complex (not self-conjugate) strategies cannot be self-tested. On the other hand, it is known that every real projective set of measurements can be embedded into a real strategy that is self-tested [CMV24]. Given the existence of self-conjugate but not real strategies, it is natural to ask whether there exist self-tests within this middle ground? In this section we give an affirmative answer to this question, and the construction arises from quaternions along with an extension of CHSH inequality.

6.1 A self-test involving quaternions

The real algebra $M_2(\mathbb{H})$ is generated by hermitian unitaries

$$h_1 = \begin{pmatrix} 1 & 0 \\ 0 & -1 \end{pmatrix}, h_2 = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}, h_3 = \begin{pmatrix} 0 & -i \\ i & 0 \end{pmatrix}, h_4 = \begin{pmatrix} 0 & -j \\ j & 0 \end{pmatrix}.$$

The elements h_1, \dots, h_4 pairwise anticommute. Under the standard $*$ -embedding $\Phi_2 : M_2(\mathbb{H}) \hookrightarrow M_4(\mathbb{C})$ from Section 2, they are represented by X_1, \dots, X_4 given as

$$\begin{aligned} X_1 &= \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & -1 & 0 \\ 0 & 0 & 0 & -1 \end{pmatrix} = \sigma_Z \otimes \text{Id}, & X_2 &= \begin{pmatrix} 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \\ 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \end{pmatrix} = \sigma_X \otimes \text{Id}, \\ X_3 &= \begin{pmatrix} 0 & 0 & -i & 0 \\ 0 & 0 & 0 & i \\ i & 0 & 0 & 0 \\ 0 & -i & 0 & 0 \end{pmatrix} = \sigma_Y \otimes \sigma_Z, & X_4 &= \begin{pmatrix} 0 & 0 & 0 & -1 \\ 0 & 0 & 1 & 0 \\ 0 & 1 & 0 & 0 \\ -1 & 0 & 0 & 0 \end{pmatrix} = \sigma_Y \otimes \sigma_Y. \end{aligned}$$

The following statement is a special case of a well-known Clifford algebra formalism [Por95].

Proposition 6.1. *If a_1, \dots, a_4 are four pairwise anticommuting hermitian unitaries on a complex Hilbert space \mathcal{H} , then the unital real subalgebra generated by them is isomorphic to $M_2(\mathbb{H})$, and there exists a unitary $U : \mathcal{H} \rightarrow \mathbb{C}^4 \otimes \mathcal{K}$ for some Hilbert space \mathcal{K} , such that*

$$U a_\ell U^* = X_\ell \otimes \text{Id} \quad \text{for } \ell = 1, \dots, 4.$$

Proof. For $\xi = (\xi_1, \xi_2, \xi_3, \xi_4) \in \mathbb{R}^4$,

$$(\xi_1 a_1 + \xi_2 a_2 + \xi_3 a_3 + \xi_4 a_4)^2 = \|\xi\|^2 1.$$

In particular, a_1, \dots, a_4 are linearly independent. Let V be the real subspace of \mathcal{A} spanned by a_1, \dots, a_4 . Then $v^2 = \|v\|^2 1$ for all $v \in V$. By the universal property of Clifford algebras [Por95, Theorem 15.13], the unital real algebra $\text{Alg}_{\mathbb{R}}(a_1, \dots, a_4)$ is a homomorphic image of the real Clifford algebra $Cl_{4,0}(\mathbb{R})$, which is isomorphic to $M_2(\mathbb{H})$ [Por95, Table 15.27]. The latter algebra is simple (i.e., has no nonzero proper ideals), so $\text{Alg}_{\mathbb{R}}(a_1, \dots, a_4) \cong M_2(\mathbb{H})$ via the map $a_\ell \mapsto h_\ell$. Finally, the complexification of $M_2(\mathbb{H})$ is the matrix algebra $\mathbb{C} \otimes_{\mathbb{R}} M_2(\mathbb{H}) \cong M_4(\mathbb{C})$, whose $*$ -embeddings into $\mathcal{B}(\mathcal{H})$ are all unitarily equivalent. Thus, \mathcal{H} factors as $\mathbb{C}^4 \otimes \mathcal{K}$ for some subspace $\mathcal{K} \subset \mathcal{H}$, and there is a unitary $U : \mathcal{H} \rightarrow \mathbb{C}^4 \otimes \mathcal{K}$ such that $U a_\ell U^* = X_\ell \otimes \text{Id}$ for $\ell = 1, \dots, 4$. \square

Let $|\phi_4\rangle = \frac{1}{2}(|00\rangle + |11\rangle + |22\rangle + |33\rangle) \in \mathbb{C}^4 \otimes \mathbb{C}^4$ be the canonical maximally entangled state of local dimension 4. The following statement can be verified by a direct calculation.

Lemma 6.2. *The largest eigenvalue of $X_1 \otimes X_1 + X_2 \otimes X_2 - X_3 \otimes X_3 + X_4 \otimes X_4$ is 4, and the corresponding eigenspace is spanned by $|\phi_4\rangle$.*

We also record a well-known fact about anticommuting unitaries.

Lemma 6.3. *Let $\theta \in \mathbb{R}$. If X and Y are anticommuting hermitian unitaries, then so are $\cos \theta X + \sin \theta Y$ and $\sin \theta X - \cos \theta Y$.*

For a bipartite state $|\psi\rangle$ and observables A_i, B_j denote

$$\text{CHSH}(|\psi\rangle; A_1, A_2; B_1, B_2) = \langle \psi | (A_1 + A_2) \otimes B_1 + (A_1 - A_2) \otimes B_2 | \psi \rangle,$$

the CHSH inequality expression. For $\ell \neq m \in \{1, \dots, 4\}$ denote

$$Y_{\ell m, \pm} = \frac{1}{\sqrt{2}}(X_\ell \pm X_m)^\Gamma.$$

Then for every pair $\ell \neq m$,

$$\text{CHSH}(|\phi_4\rangle; X_\ell, X_m; Y_{\ell m, +}, Y_{\ell m, -}) = 2\sqrt{2},$$

i.e., $(|\phi_4\rangle, \{X_\ell, X_m\}, \{Y_{\ell m, +}, Y_{\ell m, -}\})$ is an optimal strategy for the CHSH inequality. Now consider the bipartite strategy

$$S = (|\phi_4\rangle, \{X_1, X_2, X_3, X_4\}, \{Y_{\ell m, \pm} : 1 \leq \ell < m \leq 4\}). \quad (16)$$

Let J be as in Eq. (11). Since $(\text{Id}_2 \otimes J \otimes \text{Id}_2 \otimes J)|\phi_4\rangle = |\phi_4\rangle$ and the measurement algebras of S are isomorphic to $M_2(\mathbb{H})$, the strategy S is unitarily equivalent to its complex conjugate via local unitaries $\text{Id}_2 \otimes J$ (on both sides) by Eq. (12). That is, S is Schmidt self-conjugate, and non-real by Example 5.4.

Moreover, below we show that S is self-tested (in the original, real sense) by a Bell inequality (without additional assumptions on the comparing strategies; see [BCK⁺25, Theorem B.1]). To the best of our knowledge, this is the first known real self-test of a non-real strategy.

Like many other self-testing results in the literature, the following theorem exploits the well-known properties of the CHSH inequality. In particular, the Bell inequality self-testing S is a sum of six CHSH inequalities, inspired by [BŠCA18], where a sum of three CHSH inequalities is used to certify the Pauli observables. More generally, the role of Clifford algebras in identifying quantum measurements has been long recognized [Tsi87, Slo11].

Theorem 6.4. *The strategy S (16) is self-tested by the Bell inequality*

$$\sum_{1 \leq \ell < m \leq 4} \text{CHSH}(|\phi_4\rangle; X_\ell, X_m; Y_{\ell m, +}, Y_{\ell m, -}) \leq 12\sqrt{2}. \quad (17)$$

Proof. Let

$$S' = (|\psi\rangle, \{X'_1, X'_2, X'_3, X'_4\}, \{Y'_{\ell m, \pm} : 1 \leq \ell < m \leq 4\})$$

be another strategy (whose measurements are given by hermitian unitaries) on Hilbert subspaces \mathcal{H}_A and \mathcal{H}_B that attains equality in Ineq. (17) (by [BCK⁺25, Theorem B.1], it suffices to restrict to such strategies in order to establish an assumption-free self-test). Since $\text{CHSH} \leq 2\sqrt{2}$ on observables, it follows that

$$\text{CHSH}(|\psi\rangle; X'_\ell, X'_m; Y'_{\ell m, +}, Y'_{\ell m, -}) = 2\sqrt{2} \quad (18)$$

for all $\ell < m$. By the self-testing feature of the CHSH inequality (e.g., [ŠB20, Section 4]), the strategies $S_{\ell m} = (|\phi_4\rangle, \{X_\ell, X_m\}, \{Y_{\ell m, +}, Y_{\ell m, -}\})$ and $S'_{\ell m} = (|\psi\rangle, \{X'_\ell, X'_m\}, \{Y'_{\ell m, +}, Y'_{\ell m, -}\})$ give rise to the same correlation, for all $\ell < m$.

Since $\frac{1}{\sqrt{2}}(X'_\ell \pm X'_m)$ and $Y'_{\ell m, \pm}$ are hermitian unitaries (the former one by Lemma 6.3),

$$\begin{aligned} \left\| \frac{1}{\sqrt{2}}(X'_\ell \pm X'_m) \otimes \text{Id} |\psi\rangle - \text{Id} \otimes Y'_{\ell m, \pm} |\psi\rangle \right\|^2 &= 2 - 2 \langle \psi | \frac{1}{\sqrt{2}}(X'_\ell \pm X'_m) \otimes Y'_{\ell m, \pm} |\psi\rangle \\ &= 2 - 2 \langle \phi_4 | \frac{1}{\sqrt{2}}(X_\ell \pm X_m) \otimes Y_{\ell m, \pm} |\phi_4\rangle \\ &= \left\| \frac{1}{\sqrt{2}}(X_\ell \pm X_m) \otimes \text{Id} |\phi_4\rangle - \text{Id} \otimes Y_{\ell m, \pm} |\phi_4\rangle \right\|^2 \\ &= \left\| (Y_{\ell m, \pm}^\top \otimes \text{Id} - \text{Id} \otimes Y_{\ell m, \pm}) |\phi_4\rangle \right\|^2 \\ &= \frac{1}{2} \text{Tr}(Y_{\ell m, \pm} - Y_{\ell m, \pm}) = 0 \end{aligned}$$

due to $S'_{\ell m}$ and $S_{\ell m}$ having the same correlations, and the tracial property of $|\phi_4\rangle$. Thus,

$$\frac{1}{\sqrt{2}}(X'_\ell \pm X'_m) \otimes \text{Id} |\psi\rangle = \text{Id} \otimes Y'_{\ell m, \pm} |\psi\rangle \quad \text{for all } \ell < m. \quad (19)$$

In particular, the observables in S preserve the support of $|\psi\rangle$. By [BCK⁺25, Proposition 3.8 and Theorem 3.9], we can replace S with a suitable local dilation, and consequently assume that $\text{supp}_A |\psi\rangle = \mathcal{H}_A$ and $\text{supp}_B |\psi\rangle = \mathcal{H}_B$ (i.e., $|\psi\rangle$ is of full Schmidt rank).

By the self-testing feature of the CHSH inequality [ŠB20, Section 4.2] for the strategy (18), we have

$$X'_\ell X'_m + X'_m X'_\ell = 0, \quad Y'_{\ell m, +} Y'_{\ell m, -} + Y'_{\ell m, -} Y'_{\ell m, +} = 0 \quad (20)$$

for all $\ell < m$. Define $X''_\ell := \frac{1}{\sqrt{2}}(Y'_{\ell 4, \pm} + Y'_{\ell 4, -})^\top$ for $\ell = 1, 2, 3$ and $X''_4 := \frac{1}{\sqrt{2}}(Y'_{14, +} - Y'_{14, -})^\top$. Using Eq. (19), one derives that

$$Y'_{\ell m, \pm} = \frac{1}{\sqrt{2}}(X'_\ell \pm X'_m)^\top \quad \text{for all } \ell < m. \quad (21)$$

Thus, $Y'_{\ell m, \pm}$ are linear combinations of X''_1, \dots, X''_4 in the precisely the same way as $Y_{\ell m, \pm}^\top$ are linear combinations of the X_ℓ . Furthermore, X''_1, \dots, X''_4 are anticommuting hermitian unitaries by Eqs. (20) and (21), and Lemma 6.3. Consequently,

$$\text{Alg}_{\mathbb{R}}(X'_1, \dots, X'_4) \cong \text{M}_2(\mathbb{H}) \cong \text{Alg}_{\mathbb{R}}(X''_1, \dots, X''_4)$$

by Proposition 6.1. Thus, there exist unitaries $U_A : \mathcal{H}_A \rightarrow \mathbb{C}^4 \otimes \mathcal{H}_{A'}$ and $U_B : \mathcal{H}_B \rightarrow \mathbb{C}^4 \otimes \mathcal{H}_{B'}$ such that

$$U_A X'_\ell U_A^* = X_\ell \otimes \text{Id}_{\mathcal{H}_{A'}}, \quad U_B X''_\ell U_B^* = X_\ell \otimes \text{Id}_{\mathcal{H}_{B'}}, \quad \text{for } 1 \leq \ell \leq 4.$$

Furthermore,

$$U_B Y'_{\ell m, \pm} U_B^* = Y_{\ell m, \pm} \otimes \text{Id}_{\mathcal{H}_{B'}}, \quad \text{for } 1 \leq \ell < m \leq 4.$$

Finally, with a slight abuse of tensor ordering in the first line,

$$\begin{aligned} & \langle \psi | (U_A \otimes U_B)^* \left((X_1 \otimes X_1 + X_2 \otimes X_2 - X_3 \otimes X_3 + X_4 \otimes X_4) \otimes \text{Id}_{\mathcal{H}_{A'} \otimes \mathcal{H}_{B'}} \right) (U_A \otimes U_B) | \psi \rangle \\ &= \langle \psi | (X'_1 \otimes X''_1 + X'_2 \otimes X''_2 - X'_3 \otimes X''_3 + X'_4 \otimes X''_4) | \psi \rangle \\ &= \langle \phi_4 | X_1 \otimes X_1 + X_2 \otimes X_2 - X_3 \otimes X_3 + X_4 \otimes X_4 | \phi_4 \rangle \end{aligned}$$

since the correlations of $S'_{\ell m}$ and $S_{\ell m}$ coincide, and Eq. (21). Lemma 6.2 then implies that

$$(U_A \otimes U_B) | \psi \rangle = | \phi_4 \rangle \otimes | j \rangle$$

for some $| j \rangle \in \mathcal{H}_{A'} \otimes \mathcal{H}_{B'}$. Thus, S is a local dilation of S' . \square

Remark 6.5. *The self-test in Theorem 6.4 is robust: if a strategy S' attains $12\sqrt{2} - \varepsilon$ (for sufficiently small $\varepsilon > 0$) on the left-hand side of the Bell inequality (17), then S is an $\mathcal{O}(\sqrt{\varepsilon})$ -local dilation of S' [BCK⁺25, Definition 2.5]. For a wider scope of stability and rigidity results for Clifford algebras, see [Slo11, Section VI] and [RV20, Section 13.4]. Without going into technical details, let us indicate step-by-step how the robustness estimates appear in the proof of Theorem 6.4:*

- (1) *If the left-hand side of Ineq. (17) at S' is $12\sqrt{2} - \varepsilon$, then the correlations of $S_{\ell m}$ and $S'_{\ell m}$ differ up to $\mathcal{O}(\varepsilon)$, for all $\ell < m$.*
- (2) *Then, Eq. (19) holds up to $\mathcal{O}(\sqrt{\varepsilon})$; hence, after one replaces S with a strategy with a fully supported state, one may still assume that the correlations of $S_{\ell m}$ and $S'_{\ell m}$ differ up to $\mathcal{O}(\varepsilon)$ by [BCK⁺25, Lemma 3.3 and Proposition 3.4].*
- (3) *The anticommutation relations in Eq. (20) then hold on $|\psi\rangle$ up to $\mathcal{O}(\sqrt{\varepsilon})$ by [ŠB20, Section 7].*
- (4) *Next, the anticommutation characterization of $M_2(\mathbb{H})$ as in Proposition 6.1 can be also viewed through a group-theoretic lense as follows. Consider the group*

$$\Gamma = \langle g_0, \dots, g_5 \mid g_\ell^2 = 1 \text{ for } \ell \geq 0, \quad g_0 g_\ell = g_\ell g_0 \text{ for } \ell \geq 1, \quad g_\ell g_m = g_0 g_m g_\ell \text{ for } 0 < \ell < m \rangle.$$

It is easy to see that Γ is a finite group, and g_0 is central in Γ . Complex irreducible representations π of Γ are then distinguished by whether g_0 attains 1 or -1 in them. If $\pi(g_0) = 1$, then π is 1-dimensional; if $\pi(g_1) = -1$, then π is 4-dimensional and unique by Proposition 6.1. By the Gowers-Hatami theorem [GH17, Theorem 6.9] (more precisely, its version for the $|\psi\rangle$ -induced norm [Vid18]), the maps $g_\ell \mapsto X'_\ell$ and $g_\ell \mapsto X''_\ell$ are $\mathcal{O}(\sqrt{\varepsilon})$ close to compressions of representations of Γ . We may assume that these

representations are unitarily equivalent to direct powers of the 4-dimensional representation of Γ because 1-dimensional representations can be discarded if ε is small enough (as they are far away from X'_ℓ, X''_ℓ since $\text{CHSH} \leq 2$ on scalar observables). Hence, there are isometries V_A, V_B such that X_ℓ, X''_ℓ agree with $V_A^*(X_\ell \otimes \text{Id})V_A, V_B^*(X_\ell \otimes \text{Id})V_B$ on $|\psi\rangle$ up to $\mathcal{O}(\sqrt{\varepsilon})$.

- (5) Finally, a standard norm estimate on approximate eigenvectors of $(X_1 \otimes X_1 + X_2 \otimes X_2 - X_3 \otimes X_3 + X_4 \otimes X_4) \otimes \text{Id}$ and Lemma 6.2 then show that $|\psi\rangle$ agrees with $(V_A \otimes V_B)^*(|\phi_4\rangle \otimes |j\rangle)$ for some $|j\rangle$ up to $\mathcal{O}(\sqrt{\varepsilon})$.

6.2 Generating quaternions by few projections

One aspect of this paper pertains to features of strategies whose measurements generate $M_n(\mathbb{H})$ as a real algebra. In this subsection we show that small collections of projections can generate $M_n(\mathbb{H})$. While it is well known that $M_n(\mathbb{R})$ and $M_n(\mathbb{C})$ for $n \geq 2$ can be generated by three projections as real algebras (e.g., [Dav55]), the quaternion counterpart of this statement is addressed by the following proposition (which may be of independent interest in the study of real operator algebras).

Proposition 6.6. *As a real unital algebra,*

- (a) $M_n(\mathbb{H})$ for $n \in \{2, 3\}$ is generated by four projections, but not by three projections;
- (b) $M_n(\mathbb{H})$ for $n \geq 4$ is generated by three projections, but not by two projections.

Proof. First, we record the following version of Jordan's lemma: for any projections $P_1, P_2 \in M_n(\mathbb{H})$ there exists a unitary $U \in M_n(\mathbb{H})$ such that the pair (UP_1U^*, UP_2U^*) is a direct sum of pairs of 1×1 or 2×2 real matrices. Indeed, by [Zha97, Corollary 6.2], the hermitian matrix $P_1 + P_2$ can be diagonalized as $P_1 + P_2 = V\Lambda V^*$ where $V \in M_n(\mathbb{H})$ is unitary and $\Lambda \in M_n(\mathbb{R})$ is diagonal. For a column $v \in \mathbb{H}^n$ of V let $\lambda \in \mathbb{R}$ be the corresponding eigenvalue of $P_1 + P_2$. Then $\{v, P_1v\}$ spans a joint invariant subspace for P_1, P_2 on which they are represented by a pair of real matrices.

- (a) A direct calculation shows that the projections

$$\begin{pmatrix} 1 & 0 \\ 0 & 0 \end{pmatrix}, \frac{1}{2} \begin{pmatrix} 1 & 1 \\ 1 & 1 \end{pmatrix}, \frac{1}{2} \begin{pmatrix} 1 & -i \\ i & 1 \end{pmatrix}, \frac{1}{2} \begin{pmatrix} 1 & -j \\ j & 1 \end{pmatrix}$$

generate $M_2(\mathbb{H})$. Next, we show that projections $P_1, P_2 \in M_2(\mathbb{H})$ and a hermitian $Q \in M_2(\mathbb{H})$ cannot generate $M_2(\mathbb{H})$. By the first paragraph it suffices to assume that $P_1, P_2 \in M_2(\mathbb{R})$, and

$$Q = \begin{pmatrix} \alpha_1 & q \\ q^* & \alpha_2 \end{pmatrix}, \quad \alpha_\ell \in \mathbb{R}, \quad q \in \mathbb{H}.$$

By [Zha97, Lemma 2.1], there is $u \in \mathbb{H}$ of norm 1 such that $uqu^* \in \mathbb{C}$. Then the unitary $U = u\text{Id}_2 \in M_2(\mathbb{H})$ commutes with P_1 and P_2 , and $UQU^* \in M_2(\mathbb{C})$. Hence, $\text{Alg}_{\mathbb{R}}(P_1, P_2, Q)$ is isomorphic to a real subalgebra of $M_2(\mathbb{C})$, and thus distinct from $M_2(\mathbb{H})$.

Let $n = 3$. A direct calculation shows that the projections

$$\begin{pmatrix} 1 & 0 & 0 \\ 0 & 0 & 0 \\ 0 & 0 & 0 \end{pmatrix}, \frac{1}{2} \begin{pmatrix} 1 & 1 & 0 \\ 1 & 1 & 0 \\ 0 & 0 & 0 \end{pmatrix}, \frac{1}{3}v_1v_1^*, \frac{1}{3}v_2v_2^*, \quad \text{where } v_1 = \begin{pmatrix} 1 \\ i \\ 1 \end{pmatrix}, v_2 = \begin{pmatrix} 1 \\ j \\ 1 \end{pmatrix}$$

generate $M_3(\mathbb{H})$. Next, we show that projections $P_1, P_2, P_3 \in M_3(\mathbb{H})$ cannot generate $M_3(\mathbb{H})$. By the first paragraph it suffices to assume that $P_1, P_2 \in M_3(\mathbb{R})$ and $(P_1)_{13} = (P_1)_{23} = (P_2)_{13} = (P_2)_{23} = 0$. Let

$$P_3 = \begin{pmatrix} \alpha_1 & q_1 & q_2 \\ q_1^* & \alpha_2 & q_3 \\ q_2^* & q_3^* & \alpha_3 \end{pmatrix}, \quad \alpha_\ell \in \mathbb{R}, \quad q_\ell \in \mathbb{H}.$$

As in the previous paragraph, by [Zha97, Lemma 2.1] there is $u_1 \in \mathbb{H}$ of norm 1 such that $u_1 q_1 u_1^* \in \mathbb{C}$. Without loss of generality, let $q_2 \neq 0$ (the case $q_3 \neq 0$ is analogous, and the case $q_2 = q_3 = 0$ reduces to the $n = 2$ case), and denote $u_2 = \frac{u_1 q}{|u_1 q|} \in \mathbb{H}$. Then $u_1 q_2 u_2^* \in \mathbb{R}$. Since P_3 is a projection,

$$\begin{pmatrix} q_1^* & \alpha_2 & q_3 \end{pmatrix} \begin{pmatrix} \alpha_1 \\ q_1^* \\ q_2^* \end{pmatrix} = q_1^*,$$

and thus

$$(\alpha_1 + \alpha_2)(u_1 q_1^* u_1^*) + (u_1 q_3 u_2^*)(u_1 q_2 u_2^*)^* = (u_1 q_1^* u_1^*).$$

Since $u_1 q_1^* u_1^* \in \mathbb{C}$ and $u_1 q_2 u_2^* \in \mathbb{R}$, it follows that $u_1 q_3 u_2^* \in \mathbb{C}$. Then the unitary $U = u_1 \text{Id}_2 \oplus u_2 \in M_3(\mathbb{H})$ commutes with P_1 and P_2 , and $UP_3U^* \in M_3(\mathbb{C})$. Hence, $\text{Alg}_{\mathbb{R}}(P_1, P_2, P_3)$ is isomorphic to a real subalgebra of $M_3(\mathbb{C})$, and thus distinct from $M_3(\mathbb{H})$.

(b) Let $n \geq 4$. By the first paragraph of the proof, $M_n(\mathbb{H})$ cannot be generated by two projections. To see that three projections suffice, let $\theta_\ell = \frac{\pi}{\ell+2}$ for $\ell \in \mathbb{N}$, and define

$$P_1 = \bigoplus_{\ell=1}^{\lfloor n/2 \rfloor} \begin{pmatrix} 1 & 0 \\ 0 & 0 \end{pmatrix} \oplus 0^{n-2\lfloor n/2 \rfloor}, \quad P_2 = \bigoplus_{\ell=1}^{\lfloor n/2 \rfloor} \frac{1}{2} \begin{pmatrix} (\cos \theta_\ell)^2 & \cos \theta_\ell \sin \theta_\ell \\ \cos \theta_\ell \sin \theta_\ell & (\sin \theta_\ell)^2 \end{pmatrix} \oplus 0^{n-2\lfloor n/2 \rfloor}.$$

Then $P_1, P_2 \in M_n(\mathbb{R})$ are projections, and

$$\text{Alg}_{\mathbb{R}}(P_1, P_2) = \left(\bigoplus_{\ell=1}^{\lfloor n/2 \rfloor} M_2(\mathbb{R}) \right) \oplus \mathbb{R}^{n-2\lfloor n/2 \rfloor} \quad (22)$$

by, e.g., [Dav55, proof of Theorem 1]. For the third projection, we define $P_3 = \frac{1}{n} v v^* \in M_n(\mathbb{H})$ where $v \in \mathbb{H}^n$ is given as

$$v^* = (1 \quad i \quad 1 \quad j \quad 1 \quad 1 \quad \cdots \quad 1).$$

For $1 \leq \ell, m \leq n$ let $E_{\ell m}$ denote the standard matrix units (with (ℓ, m) -entry equal to one, and zeros elsewhere). By Eq. (22) we have

$$\sum_{m=1}^4 (A_m \oplus 0^{n-2}) \cdot P_3 \cdot (0^{2\ell} \oplus B_m \oplus 0^{n-2-2\ell}) \in \text{Alg}_{\mathbb{R}}(P_1, P_2, P_3) \quad (23)$$

for all $\ell \in \{0, \dots, \lfloor \frac{n}{2} \rfloor - 1\}$ and $A_m, B_m \in M_2(\mathbb{R})$. Since the first row of nP_3 is v^* , it follows by Eq. (23) that

$$iE_{11}, jE_{13}, E_{1\ell} \in \text{Alg}_{\mathbb{R}}(P_1, P_2, P_3)$$

for all $1 \leq \ell \leq n$. Since $\text{Alg}_{\mathbb{R}}(P_1, P_2, P_3)$ is closed under the conjugate transpose, we also have $E_{\ell 1} \in \text{Alg}_{\mathbb{R}}(P_1, P_2, P_3)$ for all $1 \leq \ell \leq n$, and therefore $E_{\ell m} \in \text{Alg}_{\mathbb{R}}(P_1, P_2, P_3)$ for all $1 \leq \ell, m \leq n$. Thus, $M_n(\mathbb{R}) \subset \text{Alg}_{\mathbb{R}}(P_1, P_2, P_3)$. Since we also have $iE_{11}, jE_{13} \in \text{Alg}_{\mathbb{R}}(P_1, P_2, P_3)$, it follows that $i\text{ld}, j\text{ld} \in \text{Alg}_{\mathbb{R}}(P_1, P_2, P_3)$. Thus, $\text{Alg}_{\mathbb{R}}(P_1, P_2, P_3) = M_n(\mathbb{H})$. \square

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